

# Chapter 4: Threads

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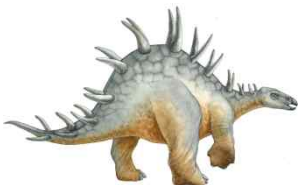
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- ❑ Overview
- ❑ Multithreading Models
- ❑ Threading Issues
- ❑ Pthreads
- ❑ Windows XP Threads
- ❑ Linux Threads
- ❑ Java Threads



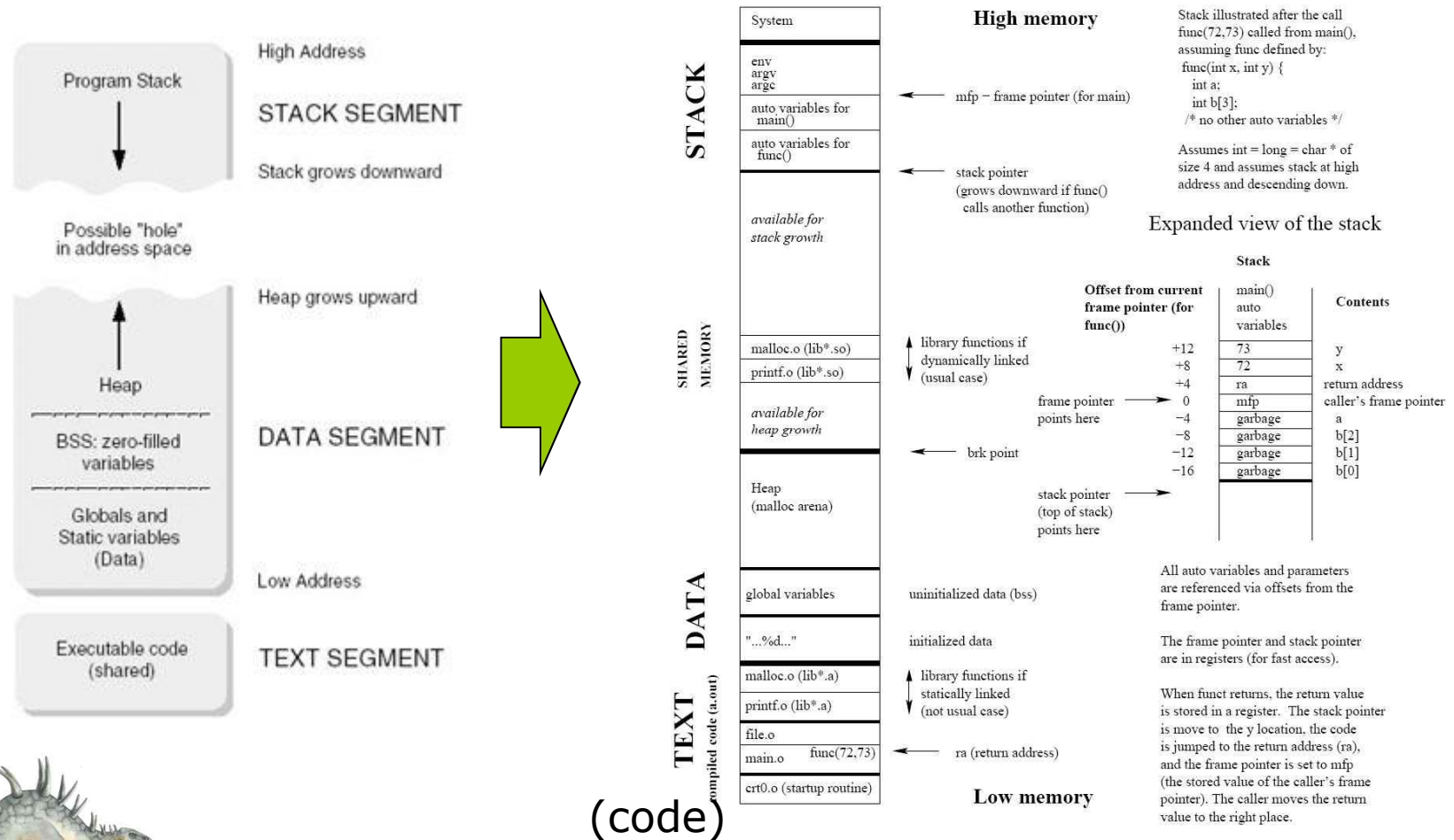
# Threads 개요

- A *thread* (or **lightweight process**) is a basic unit of CPU utilization; it consists of (보유)
  - thread ID
  - program counter
  - register set
  - stack space
  
- A thread shares with its peer threads its(공유)
  - code section
  - data section
  - operating-system resources( files ... )collectively known as a *task*.
  
- 프로세스 : 중량 프로세스(HWP;Heavy Weight Process)
  - 하나의 스레드를 가진 작업(task)



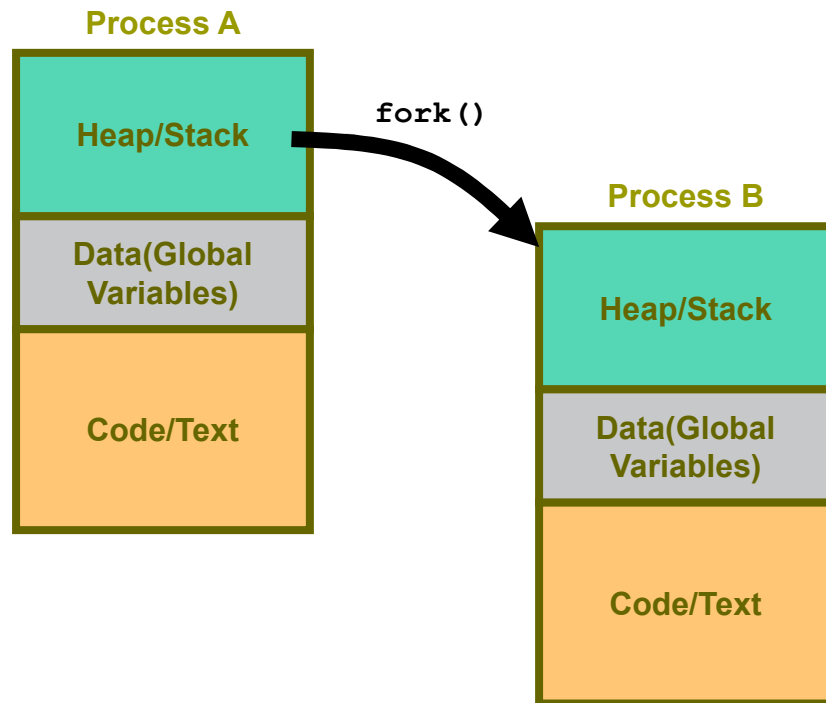
# Threads 개요

## □ Process의 메모리 구조(상세)

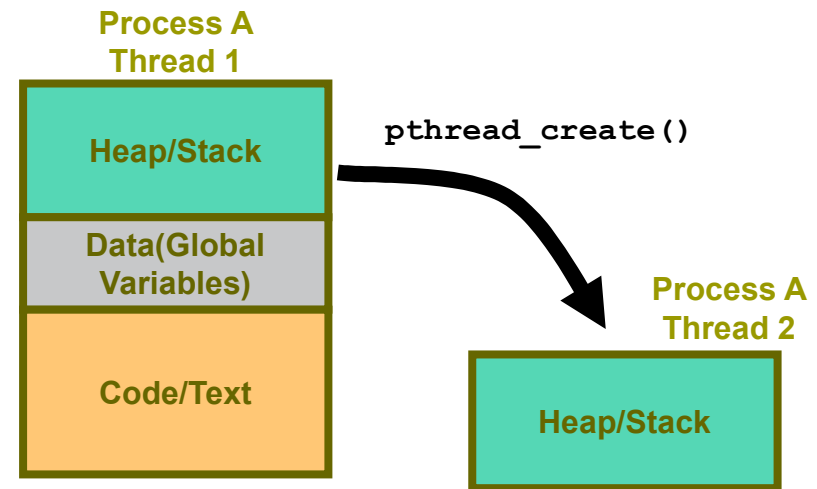
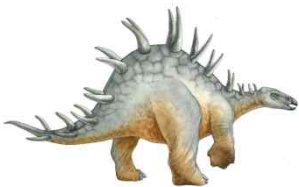


# Threads 개요

## □ Process와 Thread의 차이



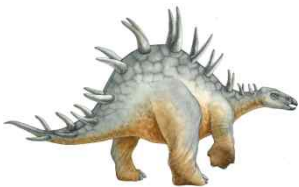
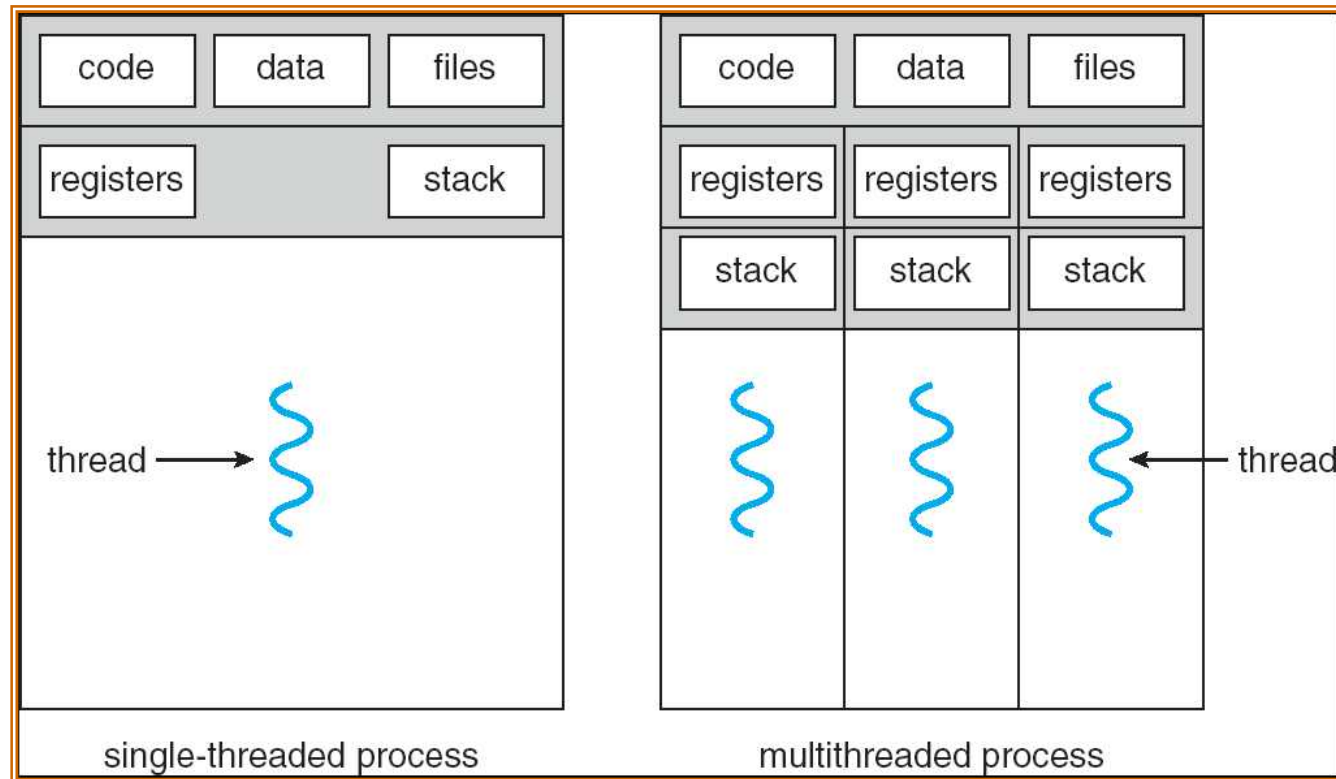
Process



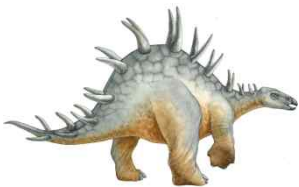
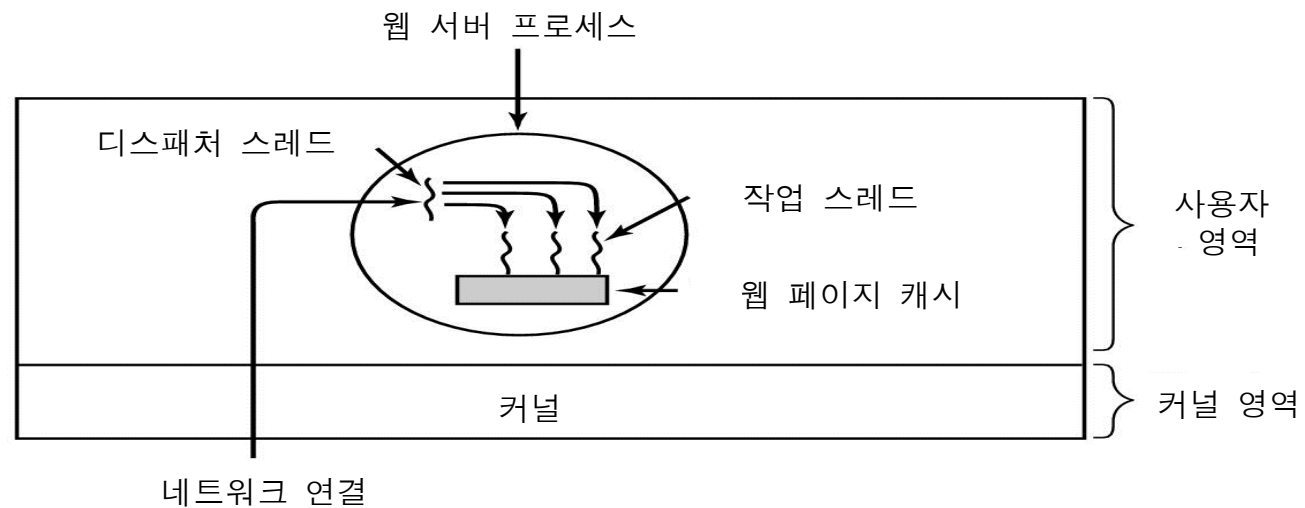
Thread



# Single and Multithreaded Processes



# 쓰레드의 이용 예 : 웹 서버



출처: 그림으로 보는 운영체제



# Benefits

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## □ Responsiveness

- eg) multi-threaded Web - if one thread is blocked (eg network) another thread continues (eg *display*)

## □ Resource Sharing

- n threads can share binary code, data, resource of the process (files, crt, ...)

## □ Economy

- creating and context switching thread (rather than a *process*)
- Solaris: 30µH 5µH

## □ Utilization of MP Architectures

- each thread may be running in *parallel* on a different processor



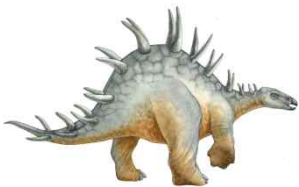


# User and Kernel Threads

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## □ User Thread

- Thread management done by user-level threads library
- 라이브러리는 커널의 지원없이 쓰레드의 생성과 스케줄링, 관리를 지원
- 커널을 통하지 않으므로, 생성과 관리가 빠르나 봉쇄형 시스템 콜을 수행하는 사용자 수준의 쓰레드는 다른 쓰레드와 함께 스케줄링 되지 않음



# User and Kernel Threads

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## □ Kernel Thread

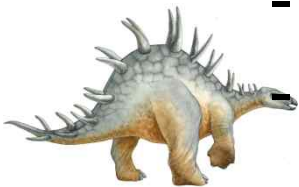
### ■ Supported by the Kernel

- 커널 수준에서 관리되어 생성과 관리가 느리나 다른 쓰레드와 함께 스케줄링 될 수 있음

### ■ Examples

- Windows 95/98/NT/2000
- Solaris
- Tru64 UNIX
- BeOS
- Linux

Java는 JVM에 의해 지원되므로,  
커널 쓰레드와 사용자 쓰레드의 중간 형태



# User and Kernel Threads

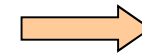
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- Some are supported by *kernel*

eg) Windows 95/98/NT

Solaris

Digital UNIX



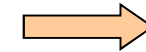
*Kernel  
Threads*

- Others are supported by *library*

eg) POSIX *Pthreads*

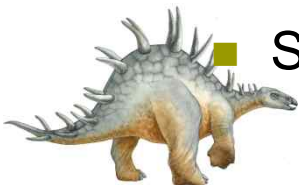
Mach *C-threads*

Solaris *threads*



*User  
Threads*

- Some are real-time threads

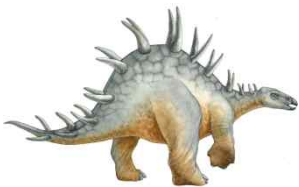


# Multithreading Models

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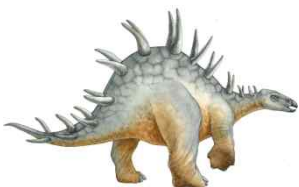
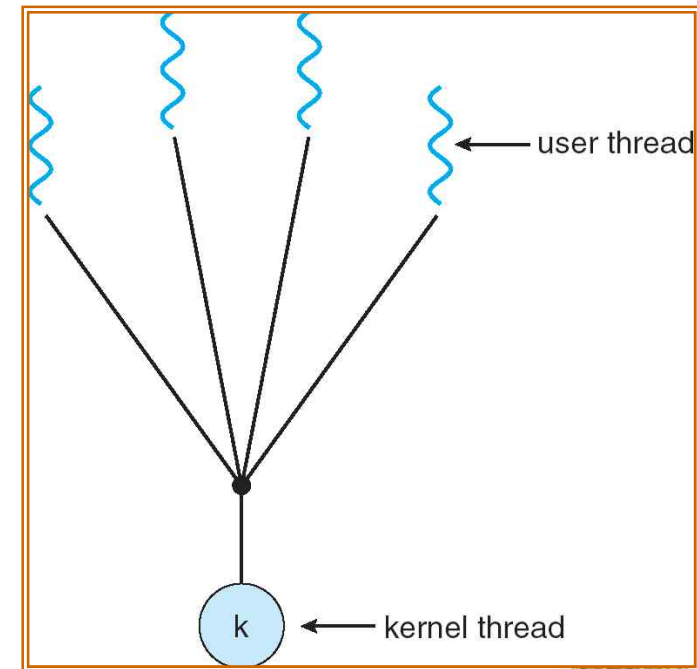
Mapping user threads to kernel threads:

- Many-to-One
- One-to-One
- Many-to-Many
- Two-level Model : Many-to-Many 모델의 변형



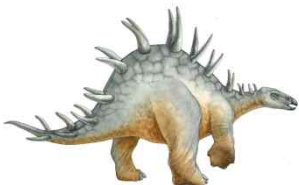
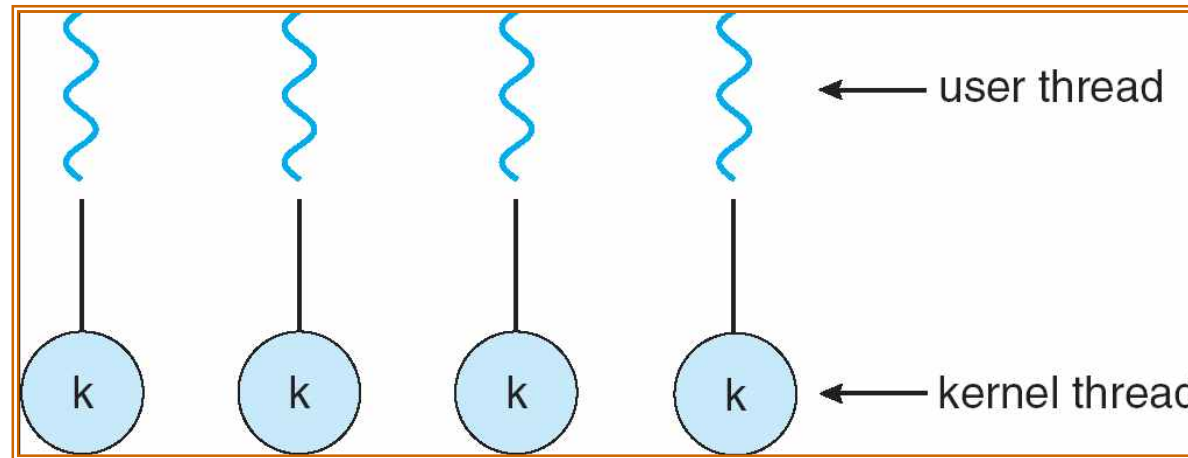
# Many-to-One

- ❑ Many user-level threads mapped to single kernel thread
- ❑ Examples:
  - Solaris Green Threads
  - GNU Portable Threads



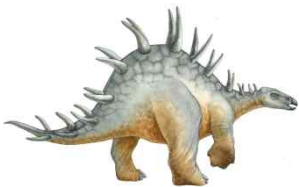
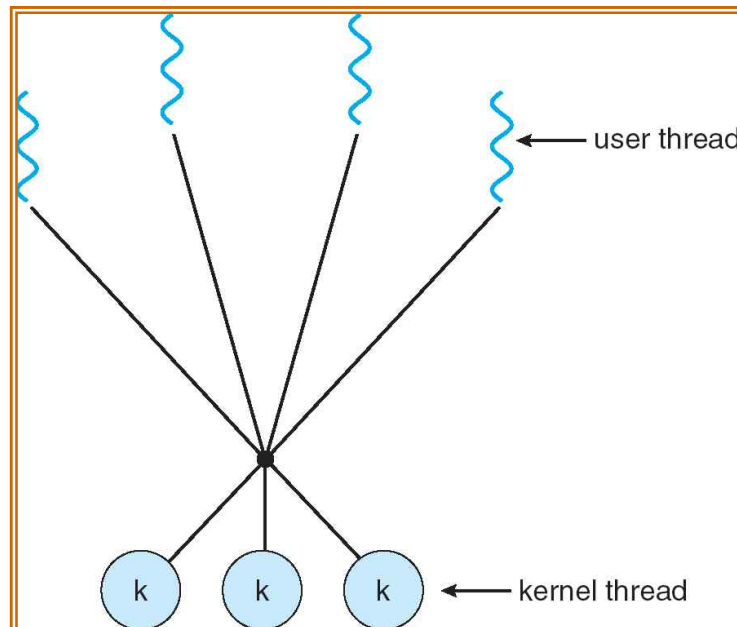
# One-to-One

- Each user-level thread maps to kernel thread
- Examples
  - Windows NT/XP/2000
  - Linux
  - Solaris 9 and later



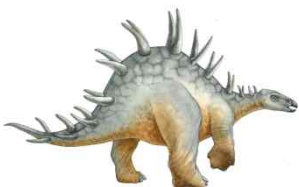
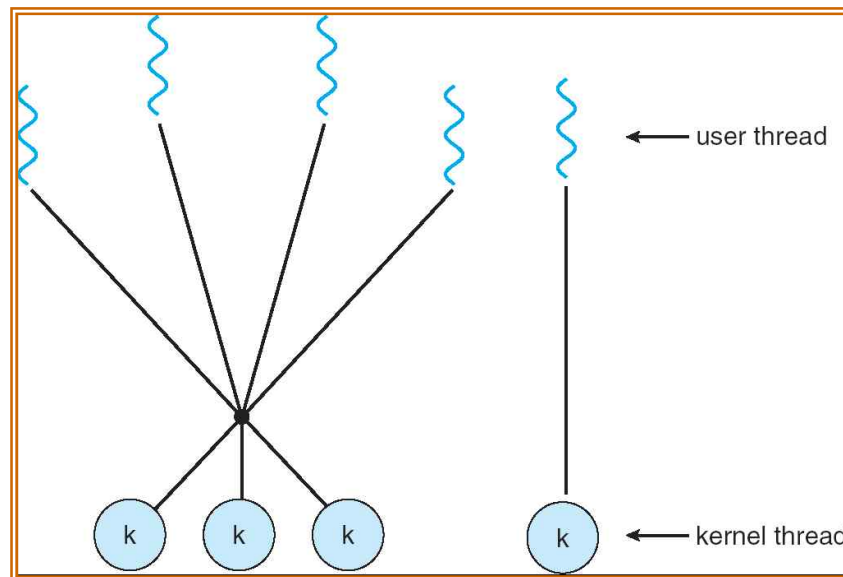
# Many-to-Many Model

- ❑ Allows many user level threads to be mapped to many kernel threads
- ❑ Allows the operating system to create a sufficient number of kernel threads
  - Solaris prior to version 9
  - Windows NT/2000 with the *ThreadFiber* package



# Two-level Model

- Similar to M:M, except that it allows a user thread to be **bound** to kernel thread
- Examples
  - IRIX, HP-UX, Tru64 UNIX, Solaris 8 and earlier

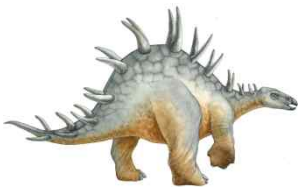




# Thread Library

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- ❑ POSIX Pthread
- ❑ Wind32 Thread API
- ❑ Java thread API
- ❑ Linux

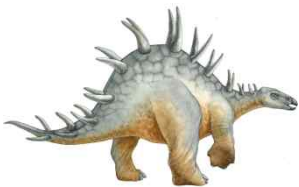


# Thread Library : Pthread

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- POSIX Pthread

- POSIX(IEEE 1003.1c)가 쓰레드 생성과 동기화를 위해 제정한 표준 API
  - Solaris, Linux, Mac OS X, Tru64 Unix에서 구현
  - 사용자 또는 Kernel 수준 라이브러리로 제공가능
  - 각 Thread는 stack의 크기와 스케줄링 정보를 가짐



# Thread Library : Pthread의 예

```
/*-----*/
/*--- main - setup server and await connections (no need to clean ---*/
/*--- up after terminated children.          ---*/
/*-----*/
int main(void)
{   int sd;
    struct sockaddr_in addr;

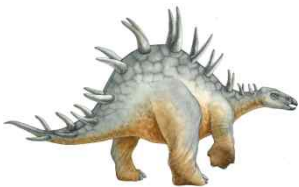
    if ( (sd = socket(PF_INET, SOCK_STREAM, 0)) < 0 )
        PANIC("Socket");
    addr.sin_family = AF_INET;
    addr.sin_port = htons(9999);
    addr.sin_addr.s_addr = INADDR_ANY;
    if ( bind(sd, (struct sockaddr*)&addr, sizeof(addr)) != 0 )
        PANIC("Bind");
    if ( listen(sd, 20) != 0 )
        PANIC("Listen");
    while (1)
    {   int client, addr_size = sizeof(addr);
        pthread_t child;

        client = accept(sd, (struct sockaddr*)&addr, &addr_size);
        printf("Connected: %s:%d\n", inet_ntoa(addr.sin_addr), ntohs(addr.sin_port));
        if ( pthread_create(&child, NULL, Child, &client) != 0 )
            perror("Thread creation");
        else
            pthread_detach(child); /* disassociate from parent */
    }
    return 0;
}
```

# Thread Library : Win32 Thread

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- Win32 Thread
  - Windows System의 Kernel 수준 라이브러리
  - Pthread 기법과 유사
    - 기본적으로 one-to-one 매핑



# Thread Library : Win32 Thread

```
#include <windows.h>
#include <iostream>

DWORD WINAPI myThread(LPVOID lpParameter)
{
    unsigned int& myCounter = *((unsigned int*)lpParameter);
    while(myCounter < 0xFFFFFFFF) ++myCounter;
    return 0;
}

int main(int argc, char* argv[])
{
    using namespace std;

    unsigned int myCounter = 0;
    DWORD myThreadID;
    HANDLE myHandle = CreateThread(0, 0, myThread, &myCounter, 0, &myThreadID);
    char myChar = ' ';
    while(myChar != 'q') {
        cout << myCounter << endl;
        myChar = getchar();
    }

    CloseHandle(myHandle);
    return 0;
}
```

The output is:

0  
868171493  
1177338657  
3782005161  
4294967295  
4294967295  
...

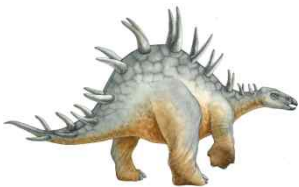


# Thread Library : Java Threads

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- Java threads are managed by the JVM
- Java threads may be created by:
  - Implementing the Runnable interface

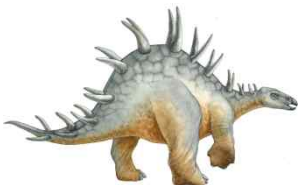
```
public interface Runnable
{
    public abstract void run();
}
```



# Java Threads - Example Program

```
class MutableInteger
{
    private int value;
    public int getValue() {
        return value;
    }
    public void setValue(int value) {
        this.value = value;
    }
}

class Summation implements Runnable
{
    private int upper;
    private MutableInteger sumValue;
    public Summation(int upper, MutableInteger sumValue) {
        this.upper = upper;
        this.sumValue = sumValue;
    }
    public void run() {
        int sum = 0;
        for (int i = 0; i <= upper; i++)
            sum += i;
        sumValue.setValue(sum);
    }
}
```



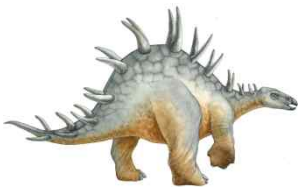
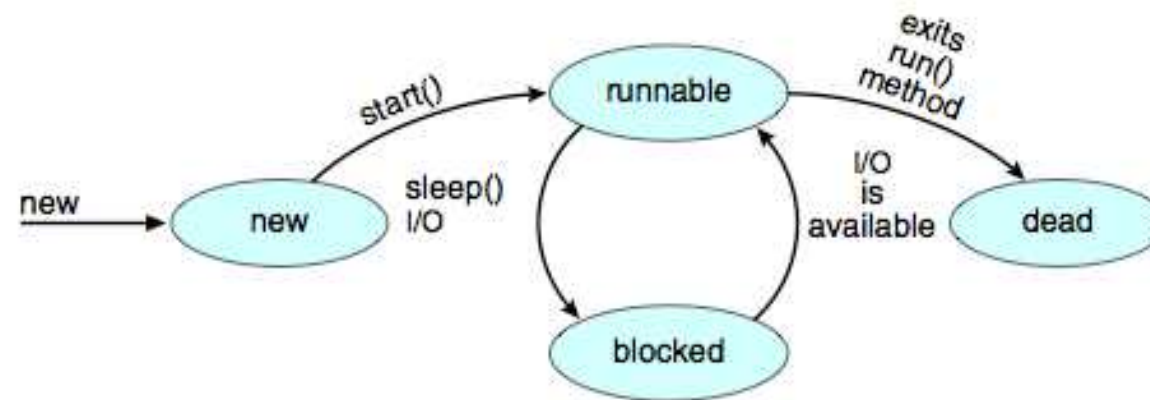
# Java Threads - Example Program

```
public class Driver
{
    public static void main(String[] args) {
        if (args.length > 0) {
            if (Integer.parseInt(args[0]) < 0)
                System.err.println(args[0] + " must be >= 0.");
            else {
                // create the object to be shared
                MutableInteger sum = new MutableInteger();
                int upper = Integer.parseInt(args[0]);
                Thread thrd = new Thread(new Summation(upper, sum));
                thrd.start();
                try {
                    thrd.join();
                    System.out.println
                        ("The sum of "+upper+" is "+sum.getValue());
                } catch (InterruptedException ie) { }
            }
        }
        else
            System.err.println("Usage: Summation <integer value>");
    }
}
```





# Java Thread States



# Java Threads - Producer-Consumer

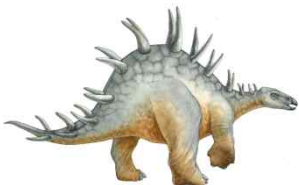
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```
public class Factory
{
    public Factory() {
        // First create the message buffer.
        Channel mailBox = new MessageQueue();

        // Create the producer and consumer threads and pass
        // each thread a reference to the mailBox object.
        Thread producerThread = new Thread(
            new Producer(mailBox));
        Thread consumerThread = new Thread(
            new Consumer(mailBox));

        // Start the threads.
        producerThread.start();
        consumerThread.start();
    }

    public static void main(String args[]) {
        Factory server = new Factory();
    }
}
```



# Java Threads - Producer-Consumer

```
class Producer implements Runnable
{
    private Channel mbox;

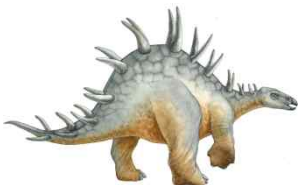
    public Producer(Channel mbox) {
        this.mbox = mbox;
    }

    public void run() {
        Date message;

        while (true) {
            // nap for awhile
            SleepUtilities.nap();

            // produce an item and enter it into the buffer
            message = new Date();

            System.out.println("Producer produced " + message);
            mbox.send(message);
        }
    }
}
```



# Java Threads - Producer-Consumer

```
class Consumer implements Runnable
{
    private Channel mbox;

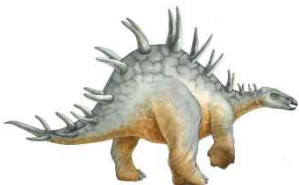
    public Consumer(Channel mbox) {
        this.mbox = mbox;
    }

    public void run() {
        Date message;

        while (true) {
            // nap for awhile
            SleepUtilities.nap();

            // consume an item from the buffer
            message = (Date)mbox.receive();

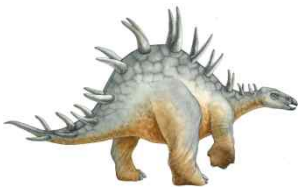
            if (message != null)
                System.out.println("Consumer consumed " + message);
        }
    }
}
```



# Threading Issues

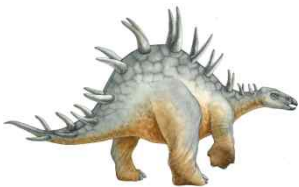
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- ❑ Semantics of **fork()** and **exec()** system calls
- ❑ Thread cancellation
- ❑ Signal handling
- ❑ Thread pools
- ❑ Thread specific data
- ❑ Scheduler activations



## Threading Issues – Semantics of fork() and exec()

- **Multithread** 프로그램에서 **fork()**를 호출한다면, 한 개의 **thread**를 생성할 것인가? 아니면 모든 **multithread**를 모두 복사해서 생성할 것인가?
- 두 개 다 지원



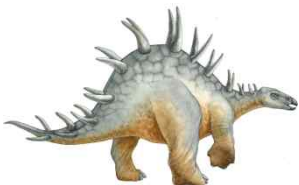
# Threading Issues – Thread Cancellation

## ❑ Terminating a thread before it has finished

- 예를 들면, 여러 스레드들이 데이터베이스를 병렬로 검색하다가 그 중 한 스레드가 결과를 찾은 경우,
- 또는 웹 브라우저에서 사용자가 **stop**을 클릭한 경우

## ❑ Two general approaches:

- **Asynchronous cancellation** terminates the target thread immediately
- **Deferred cancellation** allows the target thread to periodically check if it should be cancelled

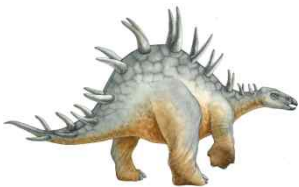


# Thread Cancellation

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Deferred cancellation in Java  
Interrupting a thread

```
Thread thrd = new Thread(new InterruptibleThread());  
thrd.start();  
.  
.  
thrd.interrupt();
```



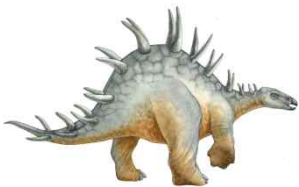


# Thread Cancellation

Deferred cancellation in Java  
Checking interruption status

```
class InterruptibleThread implements Runnable
{
    /**
     * This thread will continue to run as long
     * as it is not interrupted.
     */
    public void run() {
        while (true) {
            /**
             * do some work for awhile
             * . . .
             */

            if (Thread.currentThread().isInterrupted()) {
                System.out.println("I'm interrupted!");
                break;
            }
        }
        // clean up and terminate
    }
}
```



# Signal Handling

- Signal

- Unix에서 특정 Event가 일어났음을 알리기 위해 사용되는 단위(예: Windows Message)

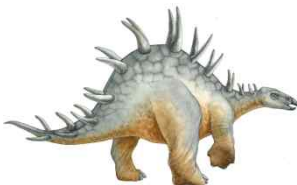
- **signal handler**의 처리 순서

1. Signal이 특정 event에 의해 생성됨
2. Signal이 특정 프로세스에 전달됨
3. Signal이 처리됨

Signal의 예  
Synchronous  
Divide-by-zero,  
illegal-memory-access

- Process에서의 Signal 처리 선택사항

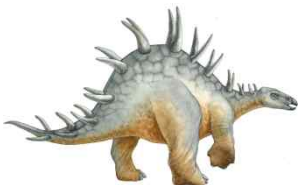
- Signal이 적용될 특정 Thread에 전송
- Process안에 있는 모든 Thread에 전송됨
- Process안의 다수 Thread에게 전송됨
- 그 Process에 전달되는 모든 Signal을 처리할 특정 Thread를 지정



# Thread Pools

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- 작업을 대기하는 다수의 Thread를 미리 생성해 놓는 Pool
- Advantages:
  - 속도 : 보통 새로운 Thread를 생성하는 것보다 존재하는 Thread를 사용하므로 다소 빠름
  - 시스템 자원 할당의 한계 설정 : Allows the number of threads in the application(s) to be bound to the size of the pool



# Thread Pools

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□ Java provides 3 thread pool architectures:

1. **Single thread executor** - pool of size 1.

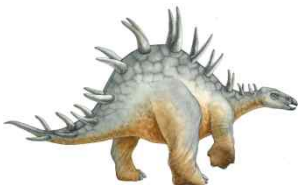
- `static ExecutorService newSingleThreadExecutor()`

2. **Fixed thread executor** - pool of fixed size.

- `static ExecutorService newFixedThreadPool(int nThreads)`

3. **Cached thread pool** - pool of unbounded size

- `static ExecutorService newCachedThreadPool()`

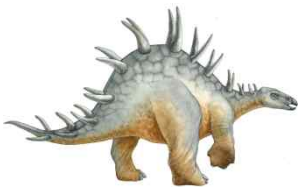


# Thread Pools

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A task to be serviced in a thread pool

```
public class Task implements Runnable
{
    public void run() {
        System.out.println("I am working on a task.");
        . . .
    }
}
```



# Thread Pools

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## Creating a thread pool in Java

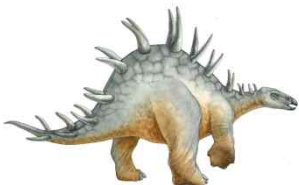
```
import java.util.concurrent.*;

public class TPEXample
{
    public static void main(String[] args) {
        int numTasks = Integer.parseInt(args[0].trim());

        // create the thread pool
        ExecutorService pool = Executors.newCachedThreadPool();

        // run each task using a thread in the pool
        for (int i = 0; i < numTasks; i++)
            pool.execute(new Task());

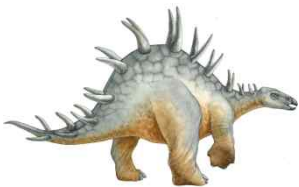
        // Shut down the pool. This shuts down the pool only
        // after all threads have completed.
        pool.shutdown();
    }
}
```



# Thread Specific Data

---

- Allows each thread to have its own copy of data
- Useful when you do not have control over the thread creation process (i.e., when using a thread pool)



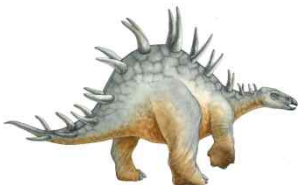
# Thread Specific Data

Thread-specific data in Java.

```
class Service
{
    private static ThreadLocal errorCode =
        new ThreadLocal();

    public static void transaction() {
        try {
            /**
             * some operation where an error may occur
             * . . .
             */
        }
        catch (Exception e) {
            errorCode.set(e);
        }
    }

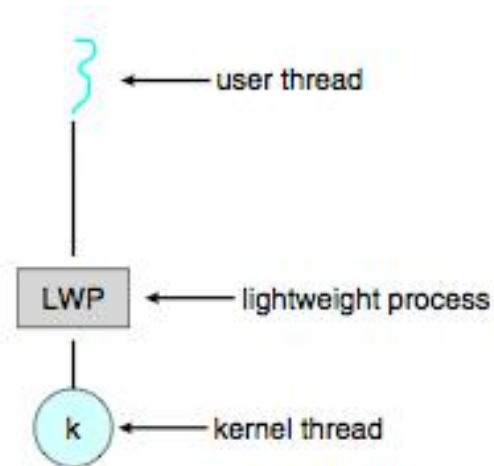
    /**
     * get the error code for this transaction
     */
    public static Object getErrorCode() {
        return errorCode.get();
    }
}
```





# Scheduler Activations

- Scheduler Activation
  - Thread library와 Kernel Thread의 통신방법
  - This communication allows an application to maintain the correct number kernel threads
- LWP 자료구조
  - M:M and Two-level model들은 다수의 Kernel



# Scheduler Activations

---

- **upcall in scheduler activation**
  - Kernel- Thread간 종료 또는 activation을 알림
    - 특정 thread 가 종료될때 upcall이 발생
    - upcall 처리기는 이 upcall을 받아 다른 thread를 activation
  - 수행되는 thread의 수를 조절



# 운영체제 사례

## □ Solaris에서 Thread와 Process의 관계

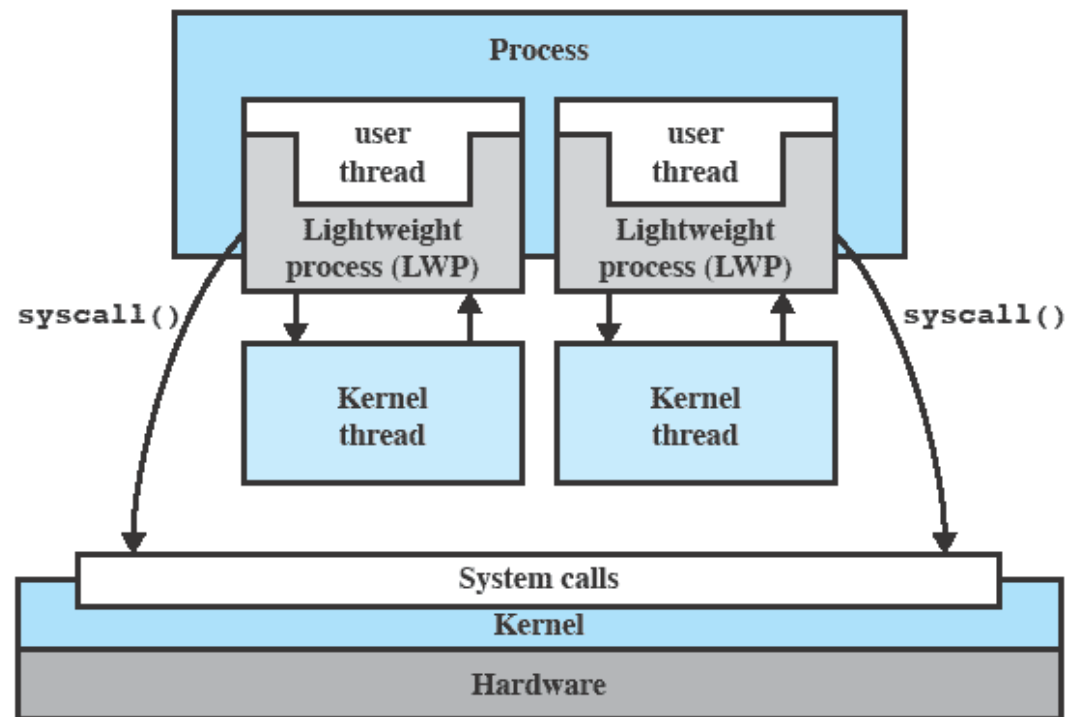


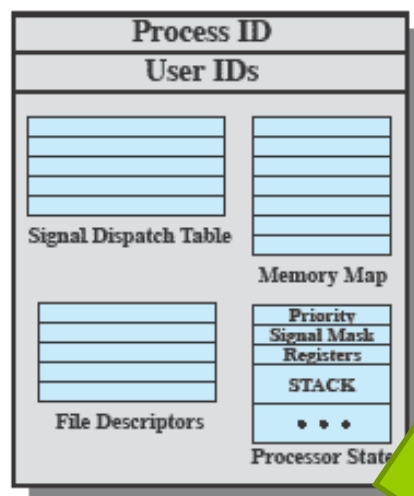
Figure 4.15 Processes and Threads in Solaris [MCDO07]



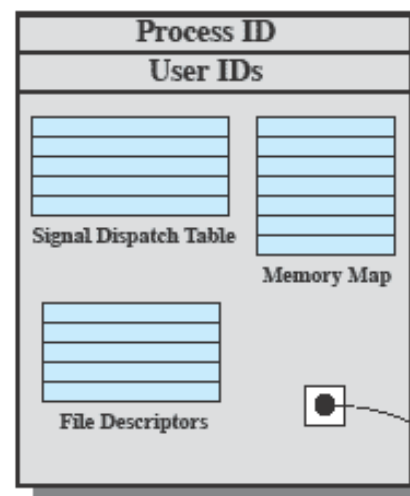
# 운영체제 사례

## □ Unix와 Solaris의 Thread 지원 Process의 비교

UNIX Process Structure



Solaris Process Structure



Solaris replaces the processor state block with a list of LWPs

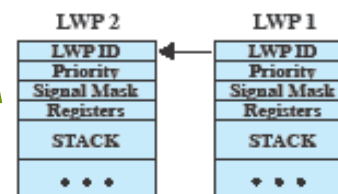


Figure 4.16 Process Structure in Traditional UNIX and Solaris [LEWI96]



# 운영체제 사례

## □ Solaris에서의 Thread 모델

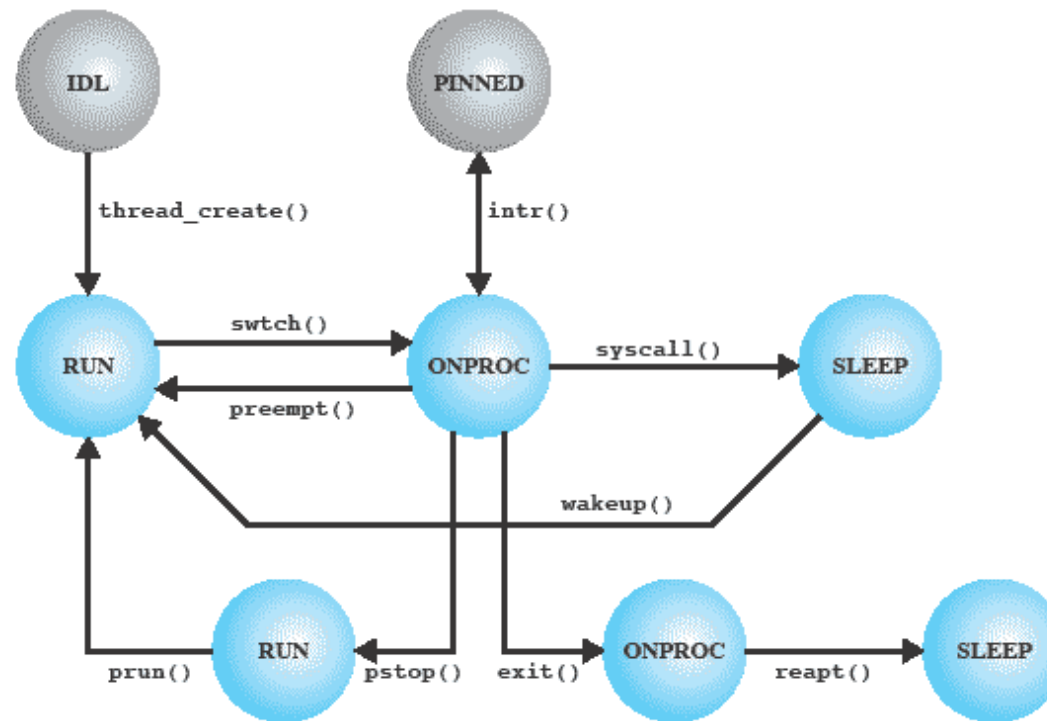


Figure 4.17 Solaris Thread States [MCDO07]



# 운영체제 사례

## □ Linux에서의 Process/Thread 모델

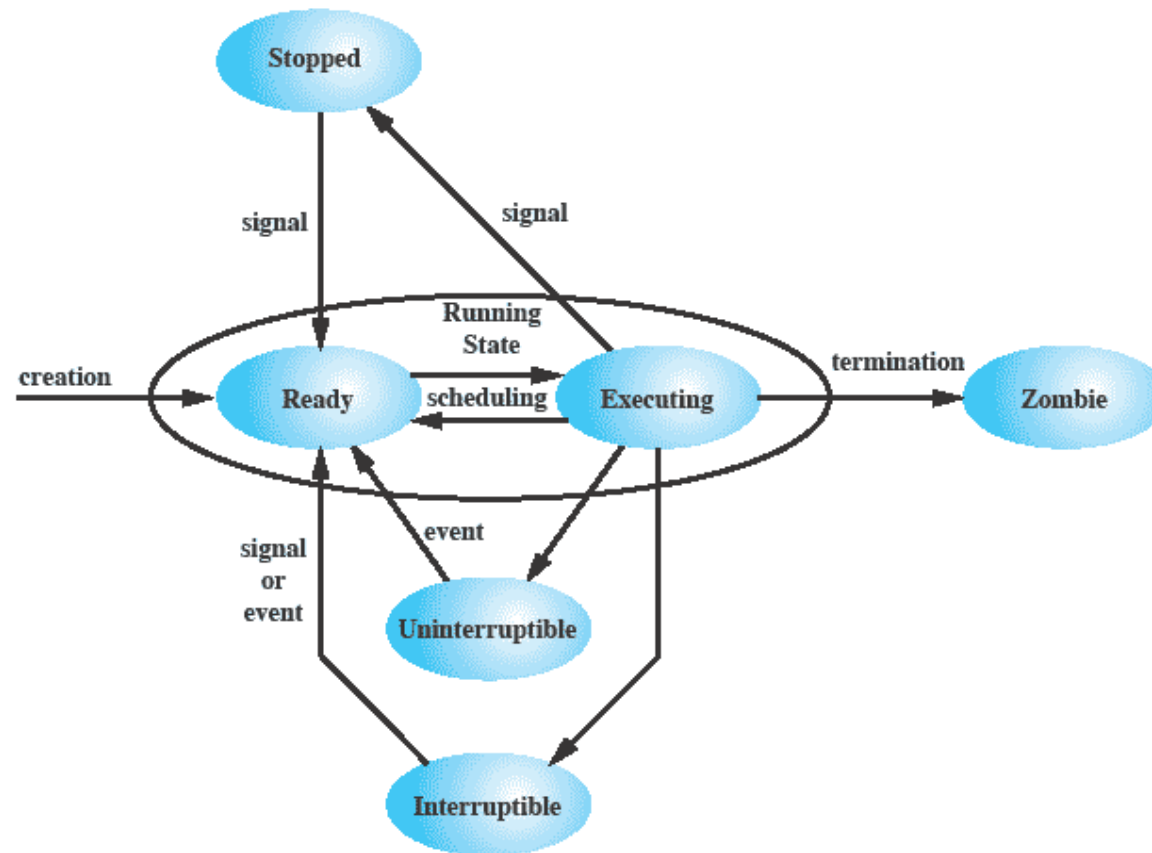


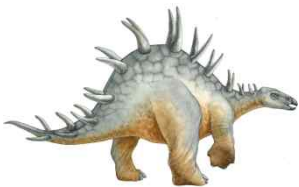
Figure 4.18 Linux Process/Thread Model



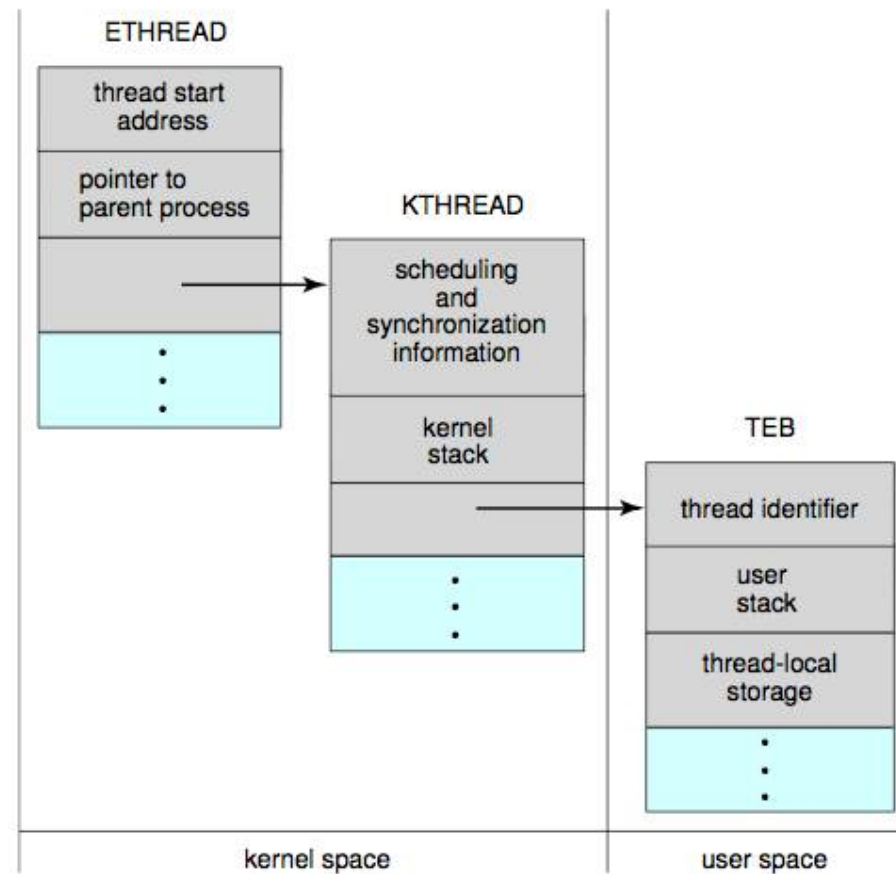
# 운영체제 사례: Windows XP Threads

---

- Implements the one-to-one mapping
- Each thread contains
  - A thread id
  - Register set
  - Separate user and kernel stacks
  - Private data storage area
- The register set, stacks, and private storage area are known as the **context** of the threads



# 운영체제 사례: Windows XP Threads

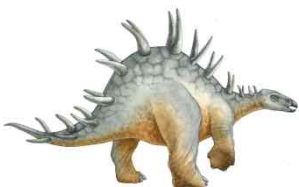




# 운영체제 사례: Linux Threads

- ❑ Linux refers to them as *tasks* rather than *threads*
- ❑ Thread creation is done through **clone()** system call
- ❑ **clone()** allows a child task to share the address space of the parent task (process)

flag	meaning
CLONE_FS	File-system information is shared.
CLONE_VM	The same memory space is shared.
CLONE_SIGHAND	Signal handlers are shared.
CLONE_FILES	The set of open files is shared.



# Thread Programming : Windows(1)

```
#include <stdio.h>
#include <string.h>
#include <windows.h>
#include <process.h>

#define rowA 3
#define colA 4
#define rowB 4
#define colB 5

typedef struct Matrix
{
    int matrixA[rowA][colA];
    int matrixB[rowB][colB];
    int matrixC[rowA][colB];
}Matrix;
unsigned long  thread0, thread1, thread2;
unsigned __stdcall Thread0(void *pParam)//스레드 함수
{
    int nTemp=0, j, k;
    Matrix *mx = (Matrix*)pParam;

    for ( j = 0; j < colB; j++ )
    {
        for ( k = 0; k < colA; k++ )
        {
            nTemp += (mx->matrixA[0][k] * mx->matrixB[k][j]);
        }
        mx->matrixC[0][j] = nTemp;
        nTemp = 0;
    }
    thread0=1;
    return 0;
}
```

# Thread Programming : Windows(2)

```
unsigned __stdcall Thread1(void *pParam)//스레드 함수
{
    int nTemp=0, j, k;
    Matrix *mx = (Matrix*)pParam;

    for ( j = 0; j < colB; j++ )
    {
        for ( k = 0; k < colA; k++ )
        {
            nTemp += (mx->matrixA[1][k] * mx->matrixB[k][j]);
        }
        mx->matrixC[1][j] = nTemp;
        nTemp = 0;
    }
    thread1=1;
    return 0;
}
```

```
unsigned __stdcall Thread2(void *pParam)//스레드 함수
{
    int nTemp=0, j, k;
    Matrix *mx = (Matrix*)pParam;

    for ( j = 0; j < colB; j++ )
    {
        for ( k = 0; k < colA; k++ )
        {
            nTemp += (mx->matrixA[2][k] * mx->matrixB[k][j]);
        }
        mx->matrixC[2][j] = nTemp;
        nTemp = 0;
    }
    thread2=1;
    return 0;
}
```

행렬곱셈

[3 \* 4] \* [4 \* 5] -> [3\*5]에서  
[1\*5] [1 \* 5] [1 \* 5]쓰레드를 통해  
[3\*5] 행렬 계산

# Thread Programming : Windows(3)

```
void main()
{
    Matrix mx;

    int i, j;
    for(i = 0; i < rowA; i++)
    {
        for(j = 0; j < colA; j++)
            mx.matrixA[i][j] = 1;
    }

    for(i = 0; i < rowB; i++)
    {
        for(j = 0; j < colB; j++)
            mx.matrixB[i][j] = 2;
    }
    _beginthreadex(NULL, 0, Thread0, &mx, 0, NULL); //스레드 시작
    _beginthreadex(NULL, 0, Thread1, &mx, 0, NULL); //스레드 시작
    _beginthreadex(NULL, 0, Thread2, &mx, 0, NULL); //스레드 시작
    while(1)
    {
        if(thread0 && thread1 && thread2)
        {
            for(i = 0; i < rowA; i++)
            {
                for(j = 0; j < colB; j++)
                    printf("%d ", mx.matrixC[i][j]);
                printf("\n");
            }
            break;
        }
    }
}
```

# 예제 : Thread Echo Server

```

/*****
*** echo-thread.c
***
*** An echo server using threads.
*****/
#include <stdlib.h>
#include <errno.h>
#include <unistd.h>
#include <string.h>
#include <sys/wait.h>
#include <sys/socket.h>
#include <resolv.h>
#include <arpa/inet.h>
#include <pthread.h>

void PANIC(char* msg);
#define PANIC(msg) { perror(msg); exit(-1); }

/*-----*/
/*--- Child - echo servlet ---*/
/*-----*/
void* Child(void* arg)
{
    char line[100];
    int bytes_read;
    int client = *(int *)arg;

    do
    {
        bytes_read = recv(client, line, sizeof(line), 0);
        send(client, line, bytes_read, 0);
    }
    while (strncmp(line, "bye\r", 4) != 0);
    close(client);
    return arg;
}

```

# 예제 : Thread Echo Server

```
/*-----*/
/*--- main - setup server and await connections (no need to clean ---*/
/*--- up after terminated children.          ---*/
/*-----*/
int main(void)
{   int sd;
    struct sockaddr_in addr;

    if ( (sd = socket(PF_INET, SOCK_STREAM, 0)) < 0 )
        PANIC("Socket");
    addr.sin_family = AF_INET;
    addr.sin_port = htons(9999);
    addr.sin_addr.s_addr = INADDR_ANY;
    if ( bind(sd, (struct sockaddr*)&addr, sizeof(addr)) != 0 )
        PANIC("Bind");
    if ( listen(sd, 20) != 0 )
        PANIC("Listen");
    while (1)
    {   int client, addr_size = sizeof(addr);
        pthread_t child;

        client = accept(sd, (struct sockaddr*)&addr, &addr_size);
        printf("Connected: %s:%d\n", inet_ntoa(addr.sin_addr), ntohs(addr.sin_port));
        if ( pthread_create(&child, NULL, Child, &client) != 0 )
            perror("Thread creation");
        else
            pthread_detach(child); /* disassociate from parent */
    }
    return 0;
}
```