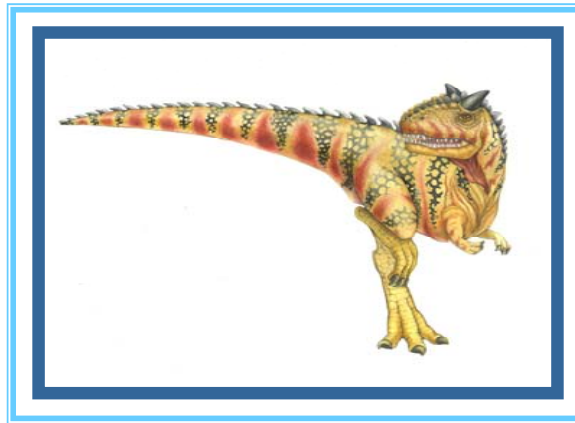


Chapter 4: Threads





Threads 개요

- A *thread* (or *lightweight process*) is a basic unit of CPU utilization; it consists of (보유)
 - thread ID
 - program counter
 - register set
 - stack space

- A thread shares with its peer threads its(공유)
 - code section
 - data section
 - operating-system resources(files ...)collectively known as a *task*.

- 프로세스 : 중량 프로세스(HWP;Heavy Weight Process)
 - 하나의 스레드를 가진 작업(task)





Motivation

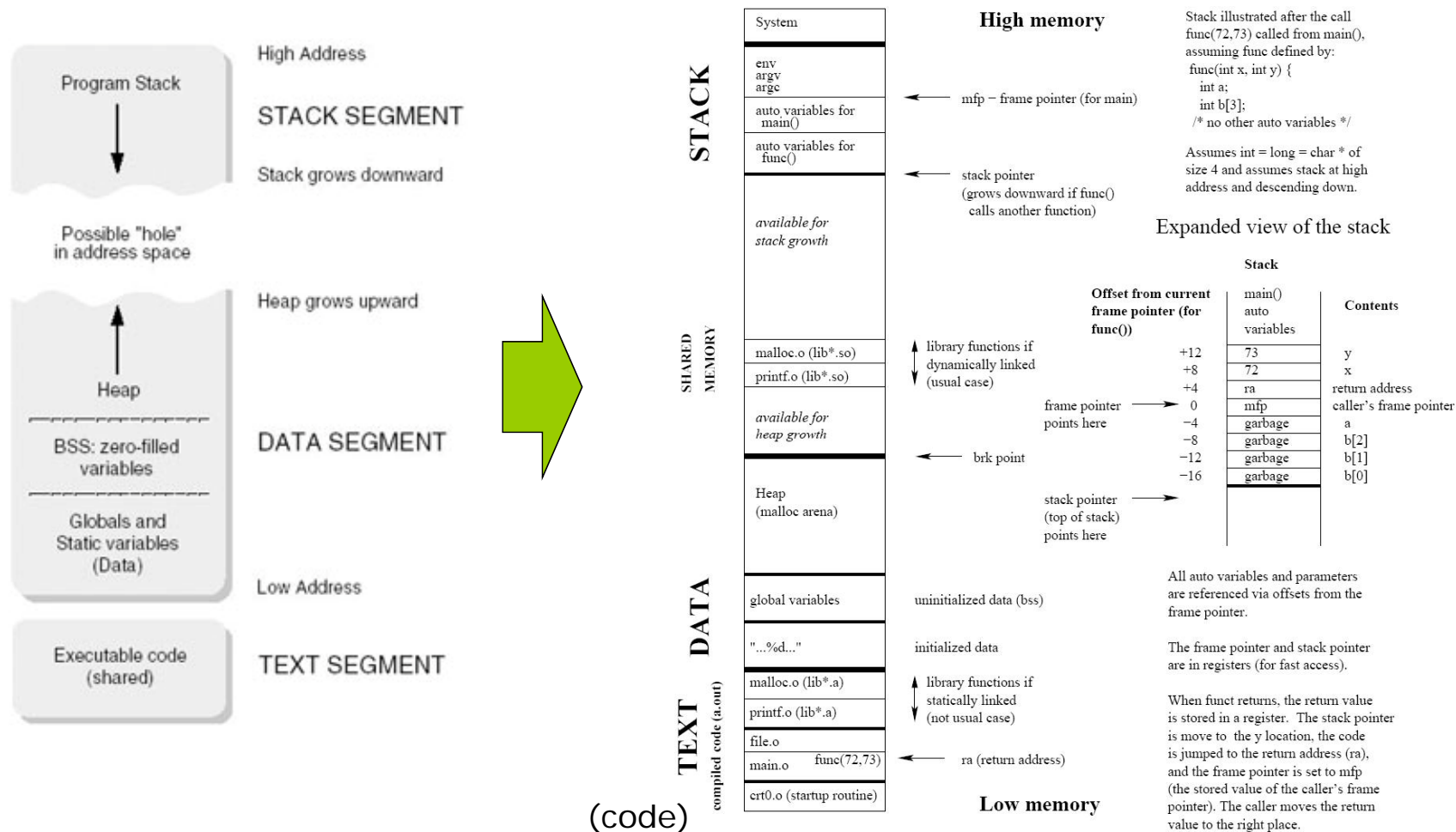
- **Thread**들은 응용내에서 수행
- 응용내에서 분리된 **thread**들로 구현되는 작업들의 예
 - **Update display**
 - **Fetch data**
 - **Spell checking**
 - **Answer a network request**
- **Kernels are generally multithreaded**





Threads 개요

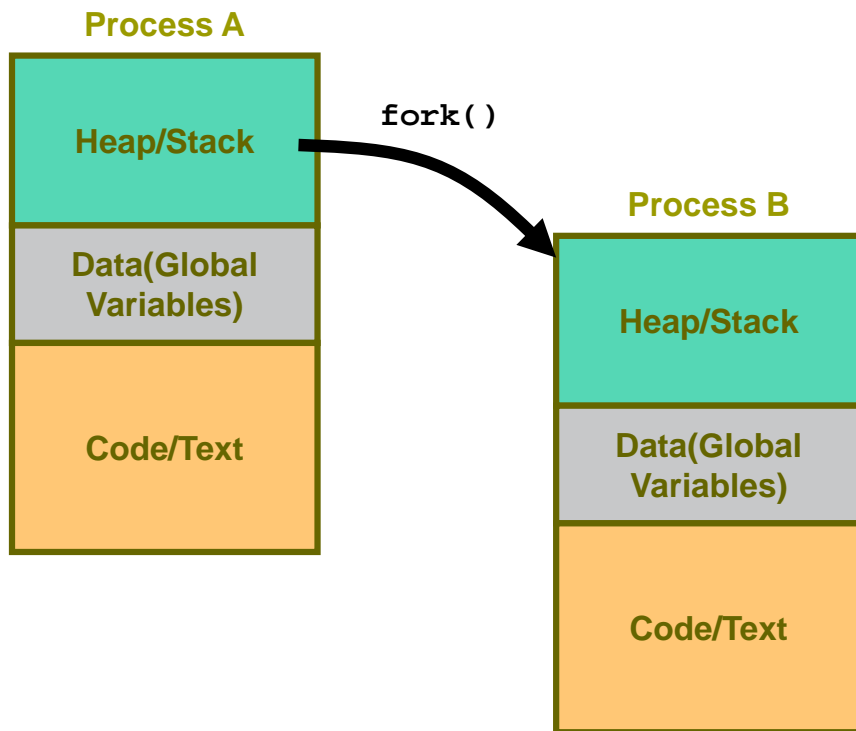
■ Process의 메모리 구조(상세)



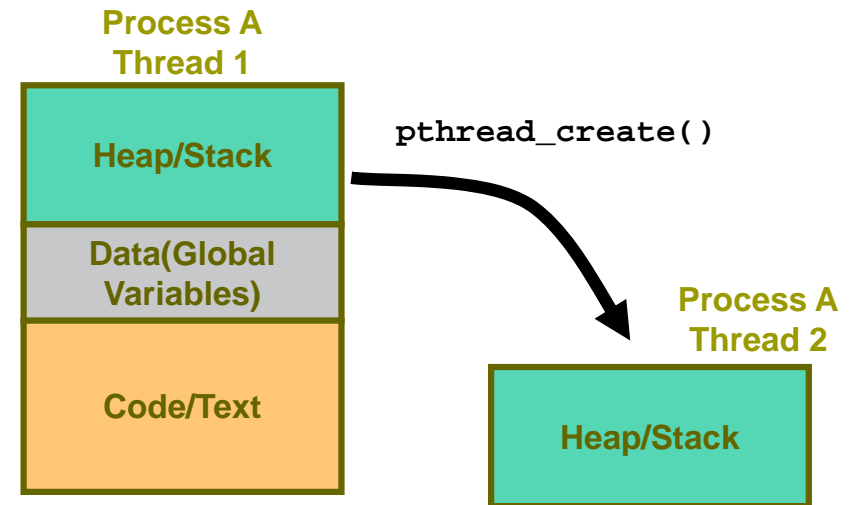


Threads 개요

■ Process와 Thread의 차이



Process

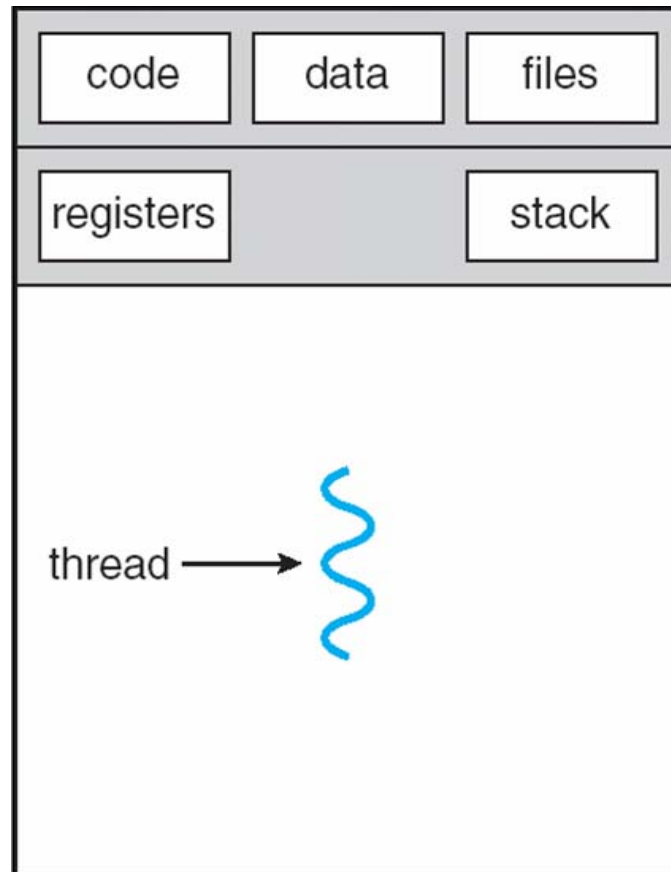


Thread

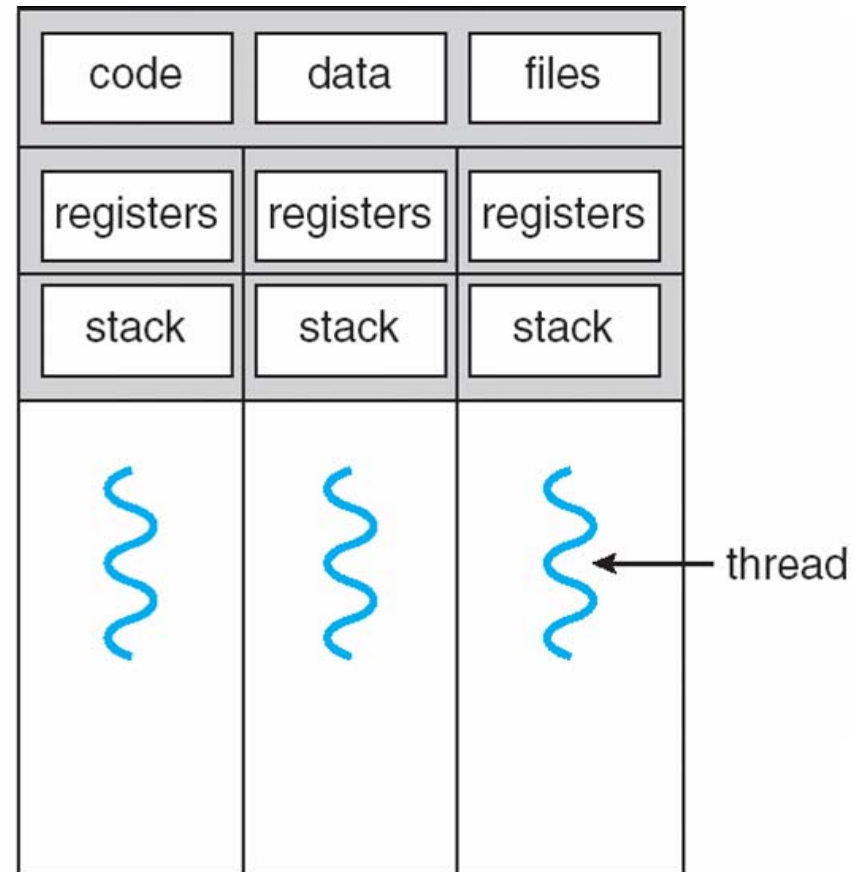




Single and Multithreaded Processes



single-threaded process



multithreaded process





Thread의 장점

■ Responsiveness

- eg) multi-threaded Web - if one thread is blocked (eg network) another thread continues (*eg display*)

■ Resource Sharing

- n threads can share binary code, data, resource of the process (files, crt, ...)

■ Economy

- creating and context switching thread (rather than a process)
- Solaris: 30배 5배

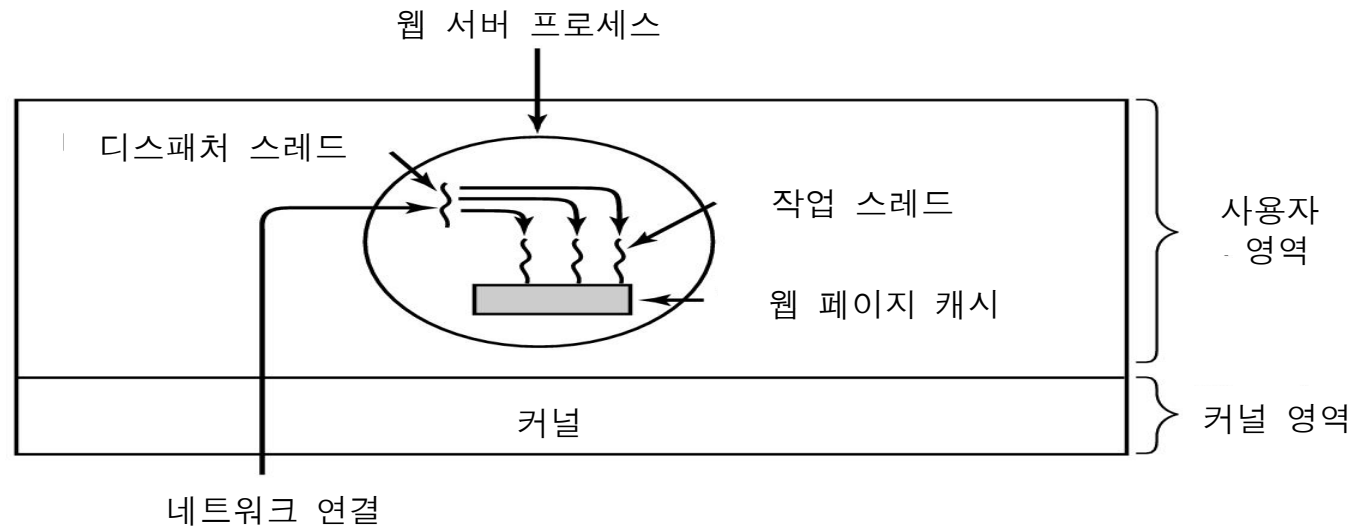
■ Utilization of MP Architectures

- each thread may be running in *parallel* on a different processor





쓰레드의 이용 예 : 웹 서버

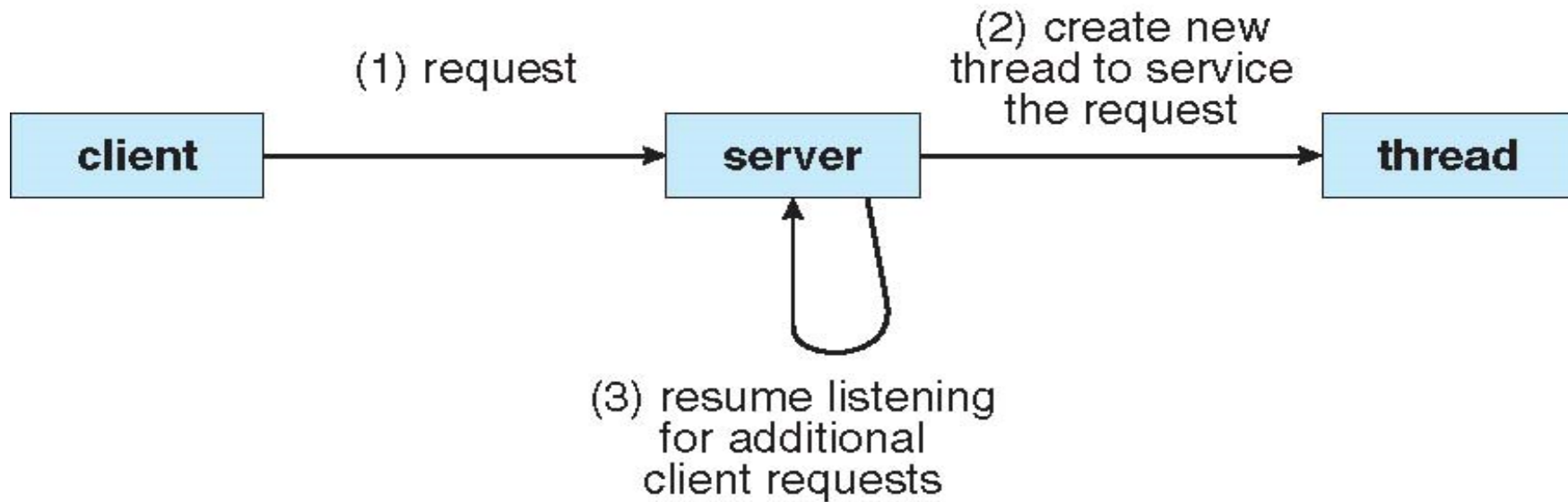


출처: 그림으로 보는 운영체제



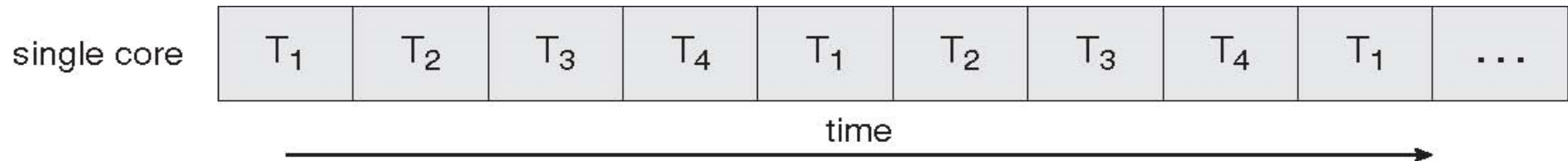


Multithreaded Server Architecture



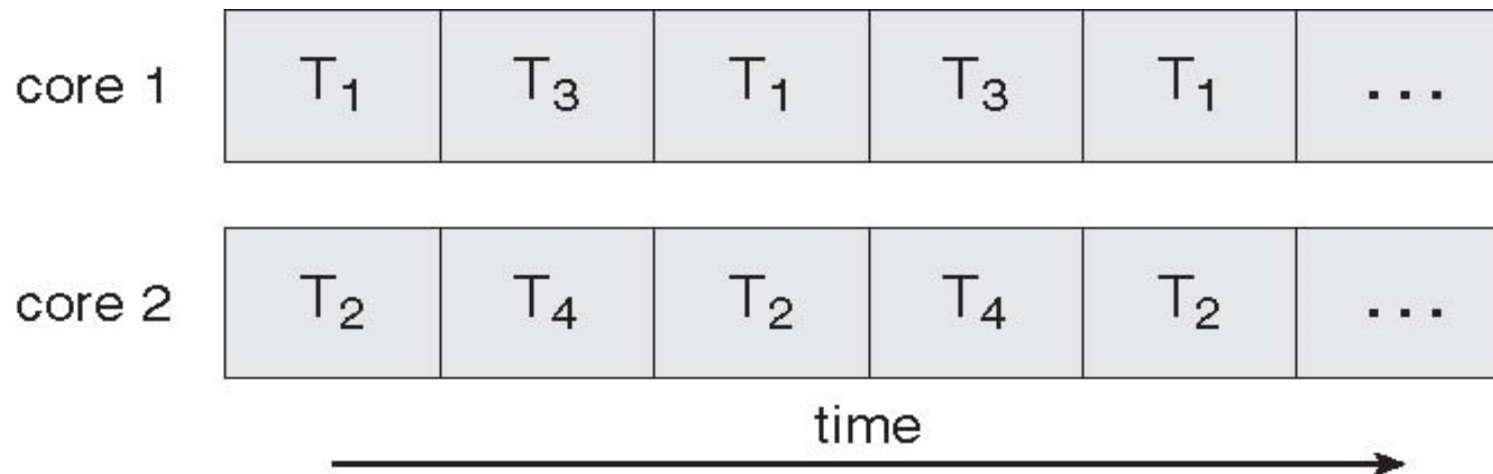


Concurrent Execution on a Single-core System





Parallel Execution on a Multicore System





User Threads

■ Thread management done by user-level threads library

- 라이브러리는 커널의 지원없이 쓰레드의 생성과 스케줄링, 관리를 지원
- 커널을 통하지 않으므로, 생성과 관리가 빠르나 봉쇄형 시스템 콜을 수행하는 사용자 수준의 쓰레드는 다른 쓰레드와 함께 스케줄링 되지 않음

■ Three primary thread libraries:

- POSIX **Pthreads**
- Win32 threads
- Java threads





Kernel Threads

- **Supported by the Kernel**
- 커널 수준에서 관리되어 생성과 관리가 느리나 다른 스레드와 함께 스케줄링 될 수 있음
- **Examples**
 - **Windows XP/2000**
 - **Solaris**
 - **Linux**
 - **Tru64 UNIX**
 - **Mac OS X**





User and Kernel Threads

- Some are supported by *kernel*

eg) Windows 95/98/NT

Solaris

Digital UNIX

→ *Kernel
Threads*

- Others are supported by *library*

eg) POSIX *Pthreads*

Mach *C-threads*

Solaris *threads*

→ *User
Threads*

- Some are real-time threads





Multithreading Models

- Many-to-One
- One-to-One
- Many-to-Many
 - Two-level Model : Many-to-Many 모델의 변형



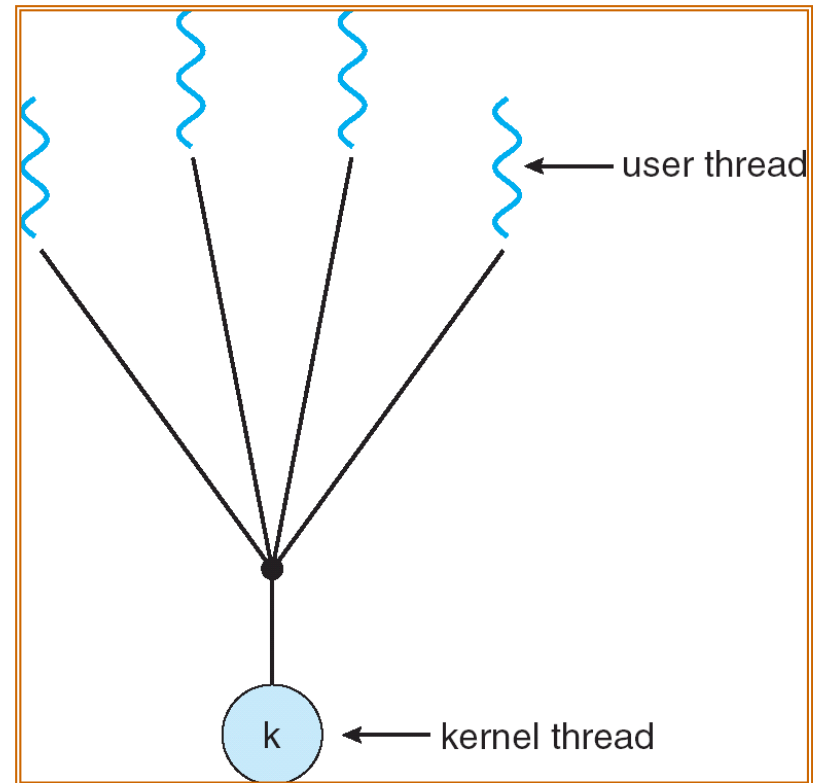


Many-to-One

- Many user-level threads mapped to single kernel thread

- Examples:

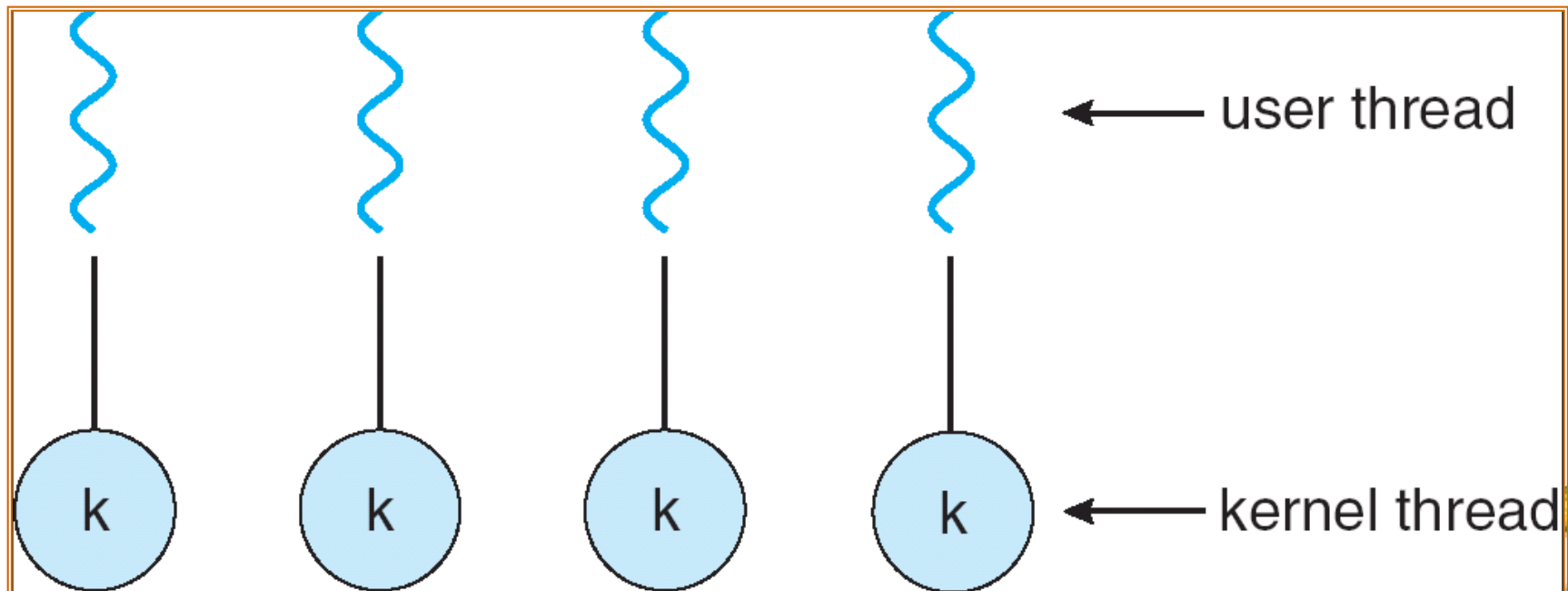
- Solaris Green Threads
- GNU Portable Threads





One-to-One

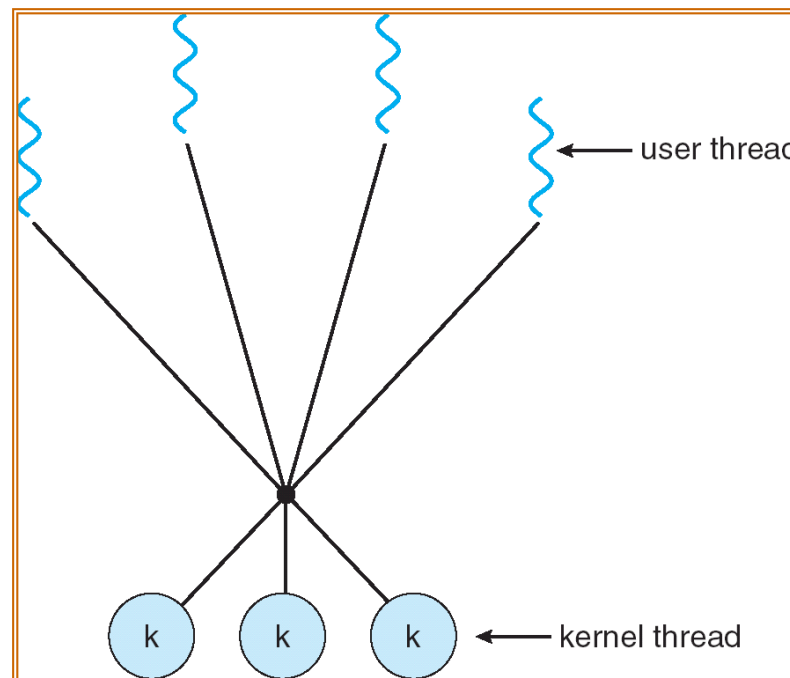
- Each user-level thread maps to kernel thread
- Examples
 - Windows NT/XP/2000
 - Linux
 - Solaris 9 and later





Many-to-Many Model

- Allows many user level threads to be mapped to many kernel threads
- Allows the operating system to create a sufficient number of kernel threads
 - Solaris prior to version 9
 - Windows NT/2000 with the *ThreadFiber* package



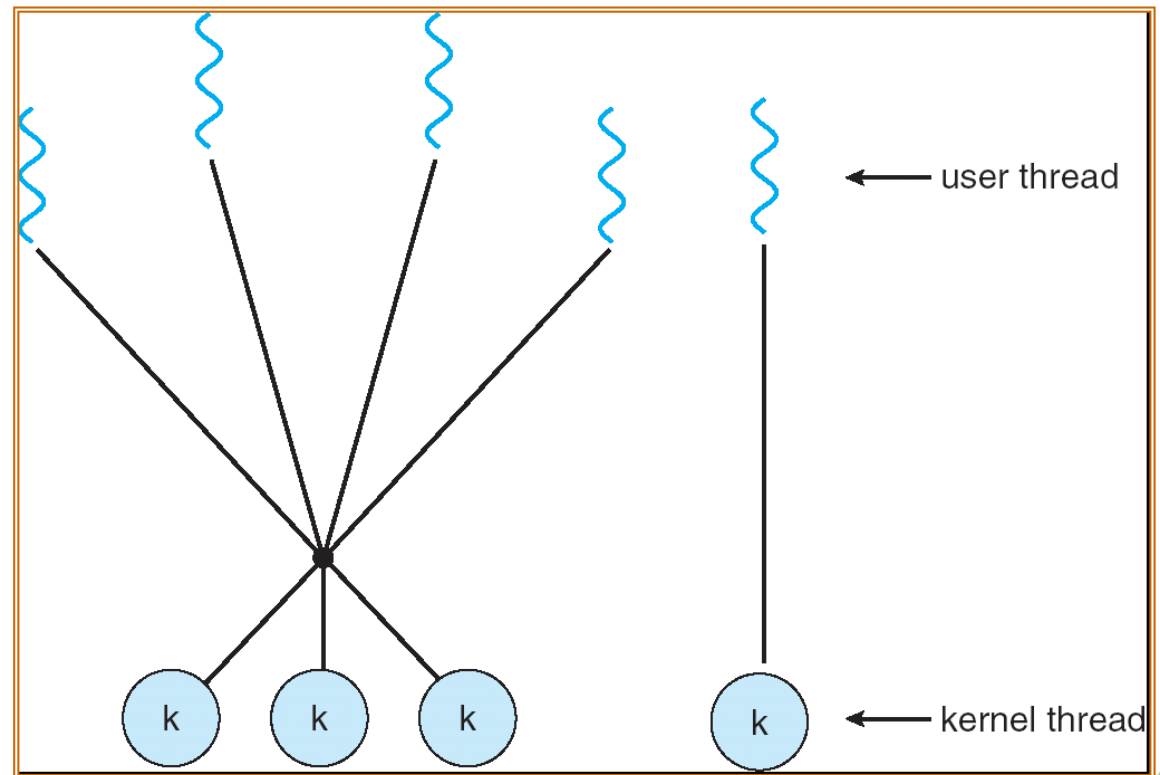


Two-level Model

- Similar to M:M, except that it allows a user thread to be bound to kernel thread

- Examples

- IRIX
- HP-UX
- Tru64 UNIX
- Solaris 8 and earlier





Thread Libraries

- **Thread library** provides programmer with API for creating and managing threads
- Two primary ways of implementing
 - Library entirely in user space
 - Kernel-level library supported by the OS
- POSIX Pthread
- Wind32 Thread API
- Java thread API
- Linux





Threading Issues

- Semantics of fork() and exec() system calls
- Thread cancellation of target thread
 - Asynchronous or deferred
- Signal handling
 - Synchronous and asynchronous





Threading Issues – Semantics of `fork()` and `exec()`

- **Multithread** 프로그램에서 **`fork()`**를 호출한다면, 한 개의 **thread**를 생성할 것인가?
아니면 모든 **multithread**를 모두 복사해서 생성할 것인가?
- 두 개 다 지원





Threading Issues – Thread Cancellation

- **Terminating a thread before it has finished**
 - 예를 들면, 여러 스레드들이 데이터베이스를 병렬로 검색하다가 그 중 한 스레드가 결과를 찾은 경우,
 - 또는 웹 브라우저에서 사용자가 **stop**을 클릭한 경우

- **Two general approaches:**
 - **Asynchronous cancellation** terminates the target thread immediately
 - **Deferred cancellation** allows the target thread to periodically check if it should be cancelled





Signal Handling

- Signals are used in UNIX systems to notify a process that a particular event has occurred.
- A **signal handler** is used to process signals
 1. Signal is generated by particular event
 2. Signal is delivered to a process
 3. Signal is handled
- Options:
 - Deliver the signal to the thread to which the signal applies
 - Deliver the signal to every thread in the process
 - Deliver the signal to certain threads in the process
 - Assign a specific thread to receive all signals for the process





Signal Handling

■ Signal

- **Unix**에서 특정 **Event**가 일어났음을 알리기 위해 사용되는 단위(예: **Windows Message**)

■ signal handler의 처리 순서

1. **Signal**이 특정 **event**에 의해 생성됨
2. **Signal**이 특정 프로세스에 전달됨
3. **Signal**이 처리됨

Signal의 예
Synchronous
Divide-by-zero,
illegal-memory-access

- **Process**에서의 **Signal** 처리 선택사항
 - ▶ **Signal**이 적용될 특정 **Thread**에 전송
 - ▶ **Process**안에 있는 모든 **Thread**에 전송됨
 - ▶ **Process**안의 다수 **Thread**에게 전송됨
 - ▶ 그 **Process**에 전달되는 모든 **Signal**을 처리할 특정 **Thread**를 지정





Thread Pools

- 작업을 대기하는 다수의 **Thread**를 미리 생성해 놓는 **Pool**

- **Advantages:**
 - 속도 : 보통 새로운 **Thread**를 생성하는 것보다 존재하는 **Thread**를 사용하므로 다소 빠름

 - 시스템 자원 할당의 한계 설정 : **Allows the number of threads in the application(s) to be bound to the size of the pool**





Thread Pools

■ Java provides 3 thread pool architectures:

1. Single thread executor - pool of size 1.

- `static ExecutorService newSingleThreadExecutor()`

2. Fixed thread executor - pool of fixed size.

- `static ExecutorService newFixedThreadPool(int nThreads)`

3. Cached thread pool - pool of unbounded size

- `static ExecutorService newCachedThreadPool()`





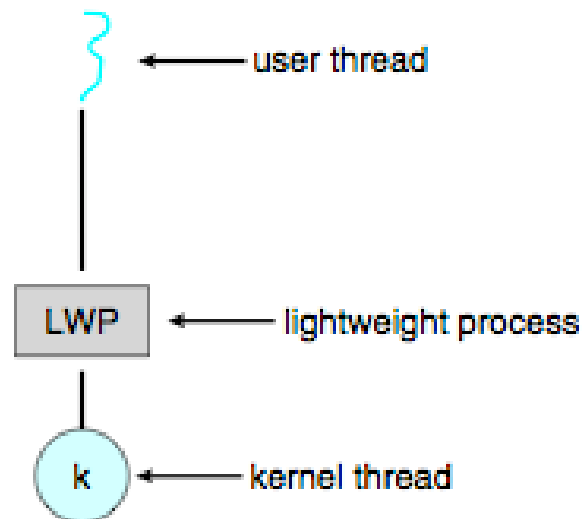
Scheduler Activations

■ Scheduler Activation

- Thread library와 Kernel Thread의 통신방법
- This communication allows an application to maintain the correct number kernel threads

■ LWP 자료구조

- M:M and Two-level model들은 다수의 Kernel





운영체제 사례

■ Solaris에서 Thread와 Process의 관계

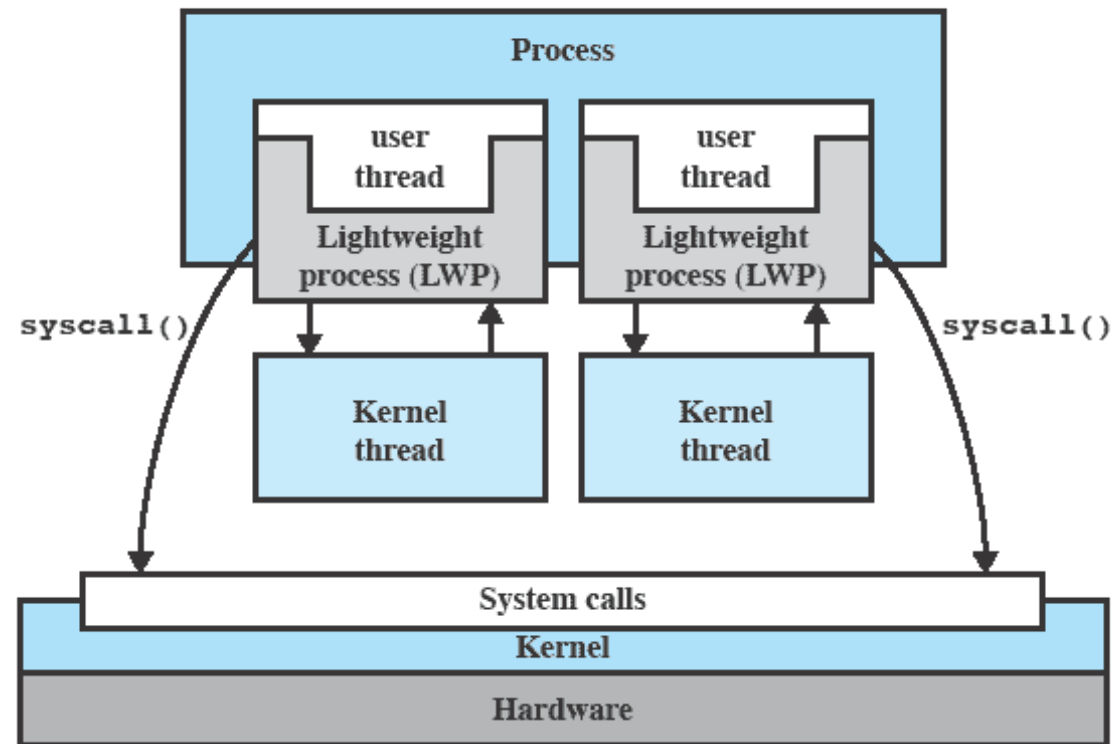


Figure 4.15 Processes and Threads in Solaris [MCDO07]





운영체제 사례

■ Unix와 Solaris의 Thread 지원 Process의 비교

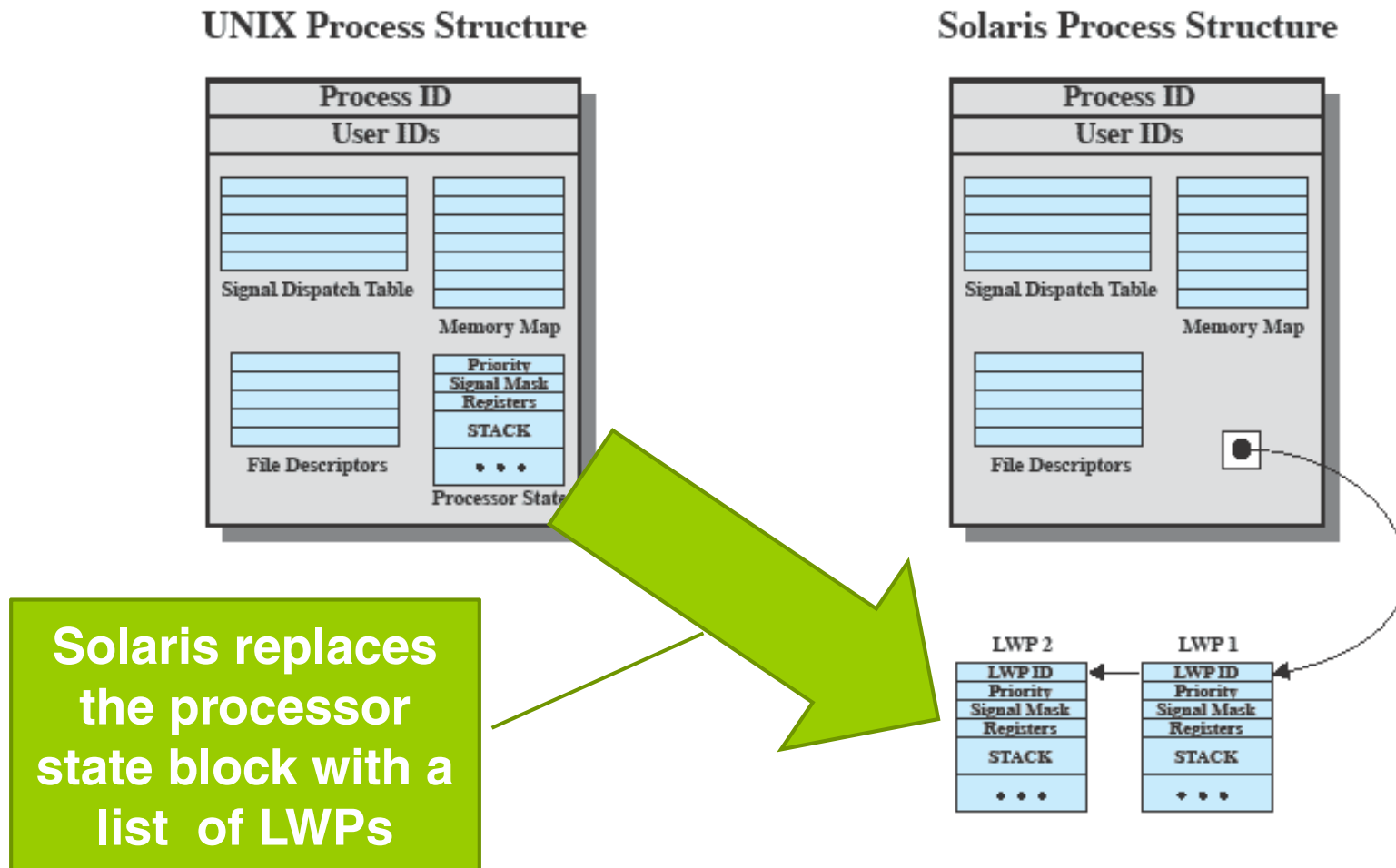


Figure 4.16 Process Structure in Traditional UNIX and Solaris [LEWI96]





운영체제 사례

■ Solaris에서의 Thread 모델

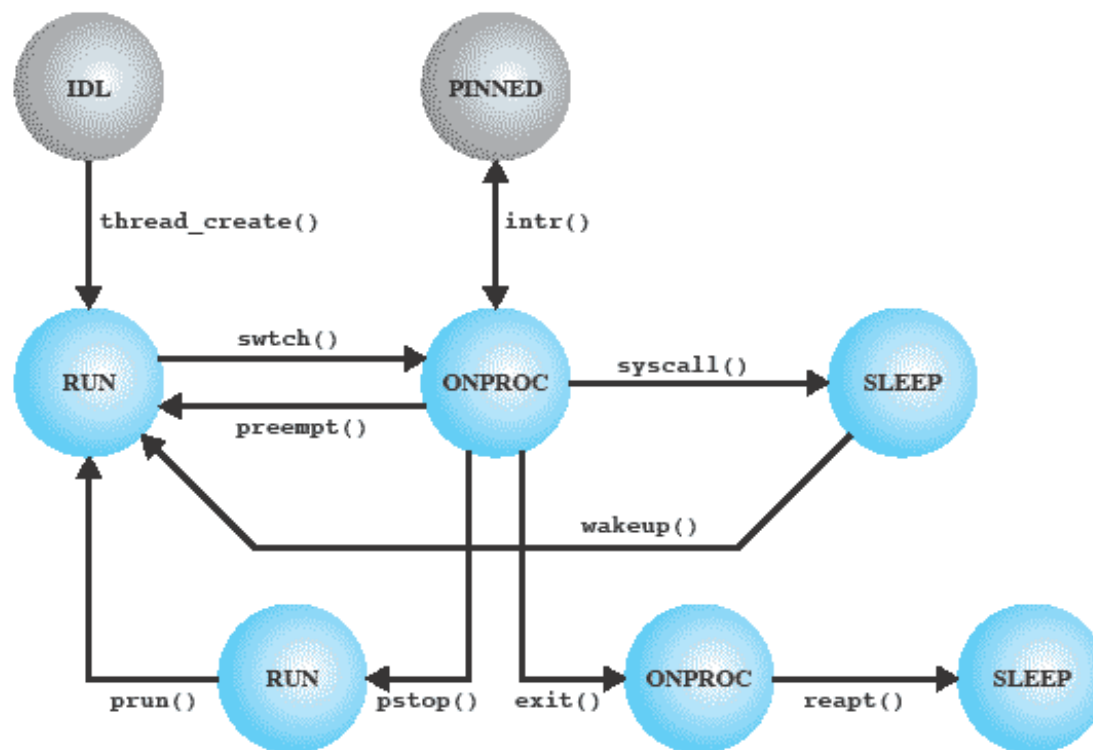


Figure 4.17 Solaris Thread States [MCDO07]





운영체제 사례

■ Linux에서의 Process/Thread 모델

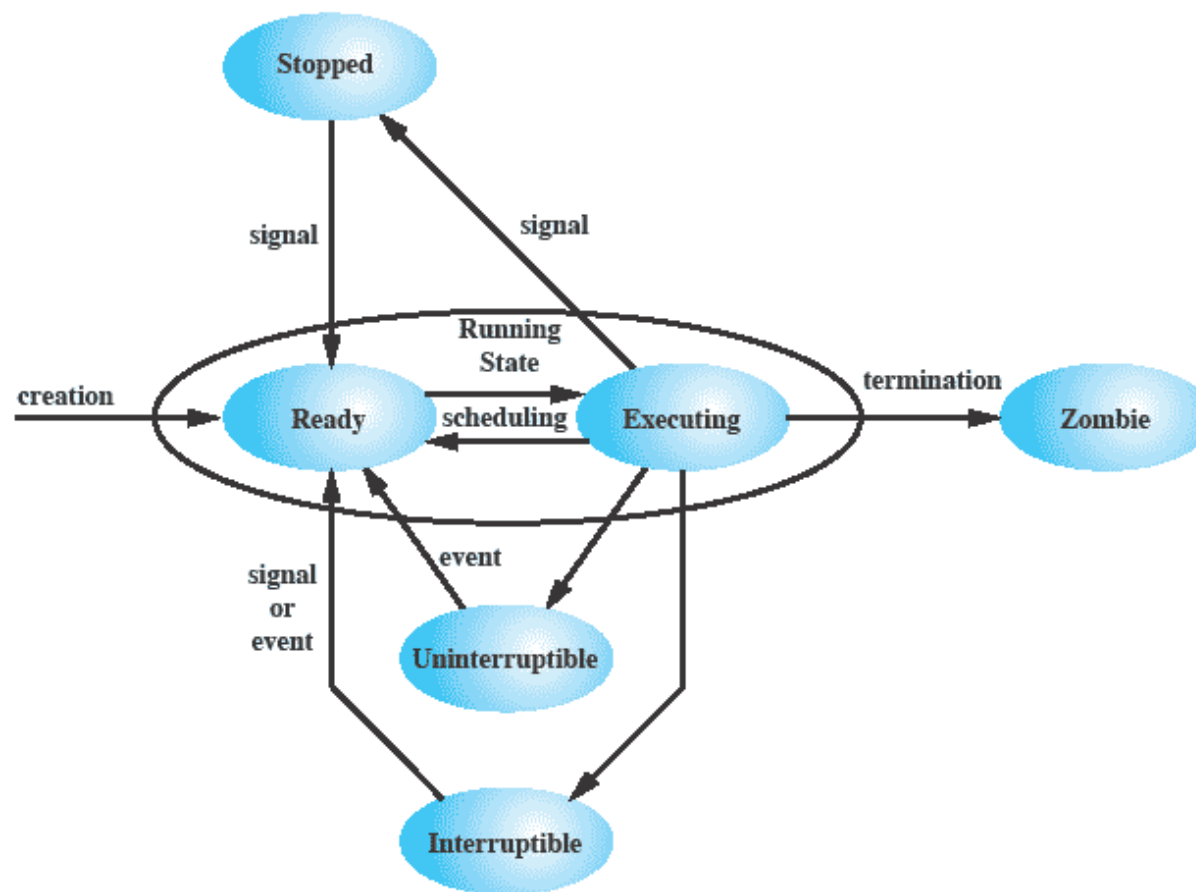


Figure 4.18 Linux Process/Thread Model





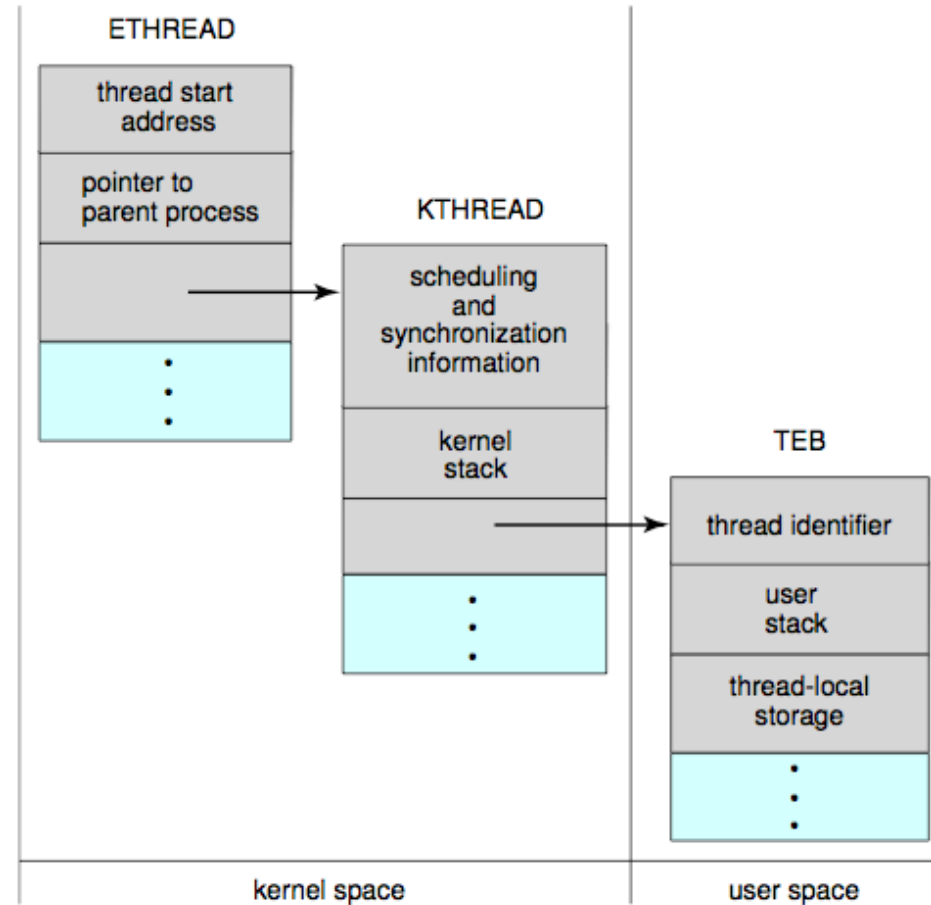
운영체제 사례: Windows XP Threads

- Implements the one-to-one mapping
- Each thread contains
 - A thread id
 - Register set
 - Separate user and kernel stacks
 - Private data storage area
- The register set, stacks, and private storage area are known as the context of the threads





운영체제 사례: Windows XP Threads





운영체제 사례: Linux Threads

- Linux refers to them as *tasks* rather than *threads*
- Thread creation is done through clone() system call
- clone() allows a child task to share the address space of the parent task (process)

flag	meaning
CLONE_FS	File-system information is shared.
CLONE_VM	The same memory space is shared.
CLONE_SIGHAND	Signal handlers are shared.
CLONE_FILES	The set of open files is shared.





Thread Programming : Windows(1)

```
#include <stdio.h>
#include <string.h>
#include <windows.h>
#include <process.h>

#define rowA 3
#define colA 4
#define rowB 4
#define colB 5

typedef struct Matrix
{
    int matrixA[rowA][colA];
    int matrixB[rowB][colB];
    int matrixC[rowA][colB];
}Matrix;

unsigned long  thread0, thread1, thread2;
unsigned __stdcall Thread0(void *pParam)//스레드 함수
{
    int nTemp=0, j, k;
    Matrix *mx = (Matrix*)pParam;

    for ( j = 0; j < colB; j++ )
    {
        for ( k = 0; k < colA; k++ )
        {
            nTemp += (mx->matrixA[0][k] * mx->matrixB[k][j]);
        }
        mx->matrixC[0][j] = nTemp;
        nTemp = 0;
    }
    thread0=1;
    return 0;
}
```





Thread Programming : Windows(2)

```
unsigned __stdcall Thread1(void *pParam)//스레드 함수
{
    int nTemp=0, j, k;
    Matrix *mx = (Matrix*)pParam;

    for ( j = 0; j < colB; j++ )
    {
        for ( k = 0; k < colA; k++ )
        {
            nTemp += (mx->matrixA[1][k] * mx->matrixB[k][j]);
        }
        mx->matrixC[1][j] = nTemp;
        nTemp = 0;
    }
    thread1=1;
    return 0;
}

unsigned __stdcall Thread2(void *pParam)//스레드 함수
{
    int nTemp=0, j, k;
    Matrix *mx = (Matrix*)pParam;

    for ( j = 0; j < colB; j++ )
    {
        for ( k = 0; k < colA; k++ )
        {
            nTemp += (mx->matrixA[2][k] * mx->matrixB[k][j]);
        }
        mx->matrixC[2][j] = nTemp;
        nTemp = 0;
    }
    thread2=1;
    return 0;
}
```

행렬곱셈

$[3 \times 4] \times [4 \times 5] \rightarrow [3 \times 5]$ 에서
 $[1 \times 5] [1 \times 5] [1 \times 5]$ 쓰레드를 통해
 $[3 \times 5]$ 행렬 계산





Thread Programming : Windows(3)

```
void main()
{
    Matrix mx;

    int i, j;
    for(i = 0; i < rowA; i++)
    {
        for(j = 0; j < colA; j++)
            mx.matrixA[i][j] = 1;
    }

    for(i = 0; i < rowB; i++)
    {
        for(j = 0; j < colB; j++)
            mx.matrixB[i][j] = 2;
    }
    _beginthreadex(NULL, 0, Thread0, &mx, 0, NULL); //스레드 시작
    _beginthreadex(NULL, 0, Thread1, &mx, 0, NULL); //스레드 시작
    _beginthreadex(NULL, 0, Thread2, &mx, 0, NULL); //스레드 시작
    while(1)
    {
        if(thread0 && thread1 && thread2)
        {
            for(i = 0; i < rowA; i++)
            {
                for(j = 0; j < colB; j++)
                    printf("%d ", mx.matrixC[i][j]);
                printf("\n");
            }
            break;
        }
    }
}
```





예제 : Thread Echo Server

```

/*****
*** echo-thread.c
***
*** An echo server using threads.
****
****/
#include <stdlib.h>
#include <errno.h>
#include <unistd.h>
#include <string.h>
#include <sys/wait.h>
#include <sys/socket.h>
#include <resolv.h>
#include <arpa/inet.h>
#include <pthread.h>

void PANIC(char* msg);
#define PANIC(msg) { perror(msg); exit(-1); }

/*-----*/
/*--- Child - echo servlet ---*/
/*-----*/
void* Child(void* arg)
{
    char line[100];
    int bytes_read;
    int client = *(int *)arg;

    do
    {
        bytes_read = recv(client, line, sizeof(line), 0);
        send(client, line, bytes_read, 0);
    }
    while (strncmp(line, "bye\r", 4) != 0);
    close(client);
    return arg;
}

```





예제 : Thread Echo Server

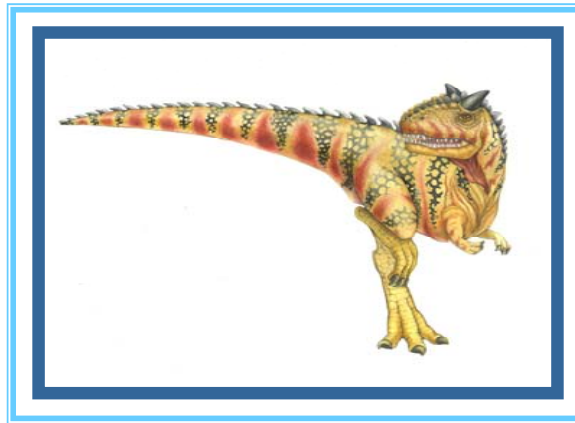
```
/*-----*/
/*--- main - setup server and await connections (no need to clean ---*/
/*--- up after terminated children. ---*/
/*-----*/
int main(void)
{   int sd;
    struct sockaddr_in addr;

    if ( (sd = socket(PF_INET, SOCK_STREAM, 0)) < 0 )
        PANIC("Socket");
    addr.sin_family = AF_INET;
    addr.sin_port = htons(9999);
    addr.sin_addr.s_addr = INADDR_ANY;
    if ( bind(sd, (struct sockaddr*)&addr, sizeof(addr)) != 0 )
        PANIC("Bind");
    if ( listen(sd, 20) != 0 )
        PANIC("Listen");
    while (1)
    {   int client, addr_size = sizeof(addr);
        pthread_t child;

        client = accept(sd, (struct sockaddr*)&addr, &addr_size);
        printf("Connected: %s:%d\n", inet_ntoa(addr.sin_addr), ntohs(addr.sin_port));
        if ( pthread_create(&child, NULL, Child, &client) != 0 )
            perror("Thread creation");
        else
            pthread_detach(child); /* disassociate from parent */
    }
    return 0;
}
```



End of Chapter 4





Pthreads

- May be provided either as user-level or kernel-level
- A POSIX standard (IEEE 1003.1c) API for thread creation and synchronization
- API specifies behavior of the thread library, implementation is up to development of the library
- Common in UNIX operating systems (Solaris, Linux, Mac OS X)





Pthreads Example

```
#include <pthread.h>
#include <stdio.h>

int sum; /* this data is shared by the thread(s) */
void *runner(void *param); /* the thread */

int main(int argc, char *argv[])
{
    pthread_t tid; /* the thread identifier */
    pthread_attr_t attr; /* set of thread attributes */

    if (argc != 2) {
        fprintf(stderr, "usage: a.out <integer value>\n");
        return -1;
    }
    if (atoi(argv[1]) < 0) {
        fprintf(stderr, "%d must be >= 0\n", atoi(argv[1]));
        return -1;
    }
}
```





Pthreads Example (Cont.)

```
/* get the default attributes */
pthread_attr_t attr;
pthread_attr_init(&attr);
/* create the thread */
pthread_create(&tid, &attr, runner, argv[1]);
/* wait for the thread to exit */
pthread_join(tid, NULL);

printf("sum = %d\n", sum);
}

/* The thread will begin control in this function */
void *runner(void *param)
{
    int i, upper = atoi(param);
    sum = 0;

    for (i = 1; i <= upper; i++)
        sum += i;

    pthread_exit(0);
}
```

Figure 4.9 Multithreaded C program using the Pthreads API.





Win32 API Multithreaded C Program

```
#include <windows.h>
#include <stdio.h>
DWORD Sum; /* data is shared by the thread(s) */
/* the thread runs in this separate function */

DWORD WINAPI Summation(LPVOID Param)
{
    DWORD Upper = *(DWORD*)Param;
    for (DWORD i = 0; i <= Upper; i++)
        Sum += i;
    return 0;
}

int main(int argc, char *argv[])
{
    DWORD ThreadId;
    HANDLE ThreadHandle;
    int Param;
    /* perform some basic error checking */
    if (argc != 2) {
        fprintf(stderr, "An integer parameter is required\n");
        return -1;
    }
    Param = atoi(argv[1]);
    if (Param < 0) {
        fprintf(stderr, "An integer >= 0 is required\n");
        return -1;
    }
}
```





Win32 API Multithreaded C Program (Cont.)

```
// create the thread
ThreadHandle = CreateThread(
    NULL, // default security attributes
    0, // default stack size
    Summation, // thread function
    &Param, // parameter to thread function
    0, // default creation flags
    &ThreadId); // returns the thread identifier

if (ThreadHandle != NULL) {
    // now wait for the thread to finish
    WaitForSingleObject(ThreadHandle, INFINITE);

    // close the thread handle
    CloseHandle(ThreadHandle);

    printf("sum = %d\n", Sum);
}
}
```

Figure 4.10 Multithreaded C program using the Win32 API.





Java Threads

- **Java threads are managed by the JVM**
- **Typically implemented using the threads model provided by underlying OS**
- **Java threads may be created by:**
 - **Extending Thread class**
 - **Implementing the Runnable interface**





Java Multithreaded Program

```
class Sum
{
    private int sum;

    public int getSum() {
        return sum;
    }

    public void setSum(int sum) {
        this.sum = sum;
    }
}

class Summation implements Runnable
{
    private int upper;
    private Sum sumValue;

    public Summation(int upper, Sum sumValue) {
        this.upper = upper;
        this.sumValue = sumValue;
    }

    public void run() {
        int sum = 0;
        for (int i = 0; i <= upper; i++)
            sum += i;
        sumValue.setSum(sum);
    }
}
```





Java Multithreaded Program (Cont.)

```
public class Driver
{
    public static void main(String[] args) {
        if (args.length > 0) {
            if (Integer.parseInt(args[0]) < 0)
                System.err.println(args[0] + " must be >= 0.");
            else {
                // create the object to be shared
                Sum sumObject = new Sum();
                int upper = Integer.parseInt(args[0]);
                Thread thrd = new Thread(new Summation(upper, sumObject));
                thrd.start();
                try {
                    thrd.join();
                    System.out.println
                        ("The sum of " + upper + " is " + sumObject.getSum());
                } catch (InterruptedException ie) { }
            }
        }
        else
            System.err.println("Usage: Summation <integer value>"); }
}
```

Figure 4.11 Java program for the summation of a non-negative integer.

